# RECOVERING MULTIPLE SMALL INCLUSIONS IN A THREE DIMENSIONAL DOMAIN USING A SINGLE MEASUREMENT

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ABSTRACT. We present and test an effective algorithm to recover multiple small heterogeneities inside arbitrary three dimensional bounded regions. In applications of electrical impedance tomography (EIT), such heterogeneities might represent tumours inside the body, weaknesses or cracks in a solid material or subterranean mineral deposits. The algorithm determines the location of all K heterogeneities using a single boundary measurement. The only nonlinear step in the algorithm is to find K roots of a single polynomial of degree K. The efficacy of this method is demonstrated by full numerical computation of the three dimensional forward problem. In addition, the methods employed can give information regarding the shape of irregular heterogeneities.

### 1. Introduction

The problem of inverse electrical impedance tomography (EIT) is to reconstruct the conductivity inside a cavity by applying a current to its periphery and measuring the response voltage. It is a classical problem which has been extensively studied since at least 1950's [1]; we refer the reader to reviews [2, 3] for an extensive list of applications and more recent works. The mathematical formulation of this problem is as follows. Let u be the voltage and  $\sigma$  be the conductivity inside a given bounded region  $\Omega \subset \mathbb{R}^d$  for d = 2, 3. An input current g(x) is applied along the boundary of the domain  $\partial\Omega$  so that

(1) 
$$\begin{cases} \nabla \cdot (\sigma \nabla u) = 0, & x \in \Omega; \\ \sigma \partial_{\nu} u = g(x), & x \in \partial \Omega, \end{cases}$$

where  $\partial_{\nu}u$  represents the outward facing normal derivative. The problem of EIT is to reconstruct the conductivity  $\sigma$  from measurements of the voltage on the boundary  $\partial\Omega$ , ie. from  $f(x) = u|_{\partial\Omega}$ .

An important special case is when the conductivity is nearly uniform except for several small defects. In applications, these defects may represent a cancer [4, 5], a buried land-mine [6], material weaknesses or corrosion in a solid structure [7, 8]. That is, consider a collection of K small non-overlapping inclusions  $\{D_j\}_{j=1}^K \subset \Omega$ , and a conductivity profile  $\sigma$  of form

(2) 
$$\sigma(x) = \begin{cases} \kappa_j, & x \in D_j, \\ 1, & x \in \Omega \setminus \bigcup_{j=1}^K D_j. \end{cases}$$

The values  $\{\kappa_j\}_{j=1}^K$  represent variable conductivities in each of the inclusions. A well-studied problem, especially in two dimensions, is to determine the location and shape of these discrete inclusions [9, 10, 11, 12, 7, 13, 14, 15, 16, 17, 18].

An effective algorithm for determining the location of a *single* defect is described in [10]. It is based on a layer potential reformulation of the problem and is applicable to arbitrary domains

in two or three dimensions. This algorithm was improved in [12] where the authors introduced a more robust variation of this method, called "the effective dipole method". These, and many other methods, use the "dipole approximation"; the fact that small inclusions effectively act as dipole forces.

For multiple inclusions, the problem was considered in [11] where an algorithm to recover K inclusions using only two boundary measurements was given. This was improved in more recent work by [9], where the authors give an algorithm which recovers multiple inclusions using a single measurement.

In the present work, we develop and test an algorithm for recovery of multiple inclusions inside an arbitrary three dimensional domain using a single measurement. The algorithm is an extension of the two-dimensional algorithm in [9] to three dimensions, using ideas developed in [19, 20, 21, 22]. In two dimensions, we also show that the algorithm developed in [12] is a special case of the method presented in [9]. The method proposed herein generalises the very recent treatment of (c.f. [23]) to the problem of finding a single dipole inside a sphere. For spherically shaped inclusions, in addition to recovering their locations, the algorithm also recovers their radius, provided that  $\kappa_j$  is known. Moreover as we illustrate on several numerical examples, irregularly shaped inclusions (such as a box) can often be effectively approximated and recovered by multiple spherical inclusions.

In our derivation of the algorithm, the well known (cf. [13, 18]) asymptotic formulation of solutions to (1) in the limit of small inclusion plays a crucial role. More specifically, if each inclusion has a common scaling  $\bar{D}_j = \varepsilon^{-1}D_j$  where  $\bar{D}_j$  are fixed, then as  $\varepsilon \to 0$ ,

(3) 
$$u(x) - u_0(x) = -\varepsilon^d \sum_{j=1}^K \nabla u_0(\xi_j) \, \mathcal{M}_j \, \nabla_y N(x, \xi_j) + \mathcal{O}(\varepsilon^{d+1}).$$

holds at points  $x \in \Omega$  away from the inclusion centres  $\{\xi_j\}_{j=1}^K$ . An associated transmission problem can be formulated for the voltage u in the case where  $|x - \xi_j| \ll 1$  [3], however, in the present work we only evaluate  $u(x) - u_0(x)$  along  $\partial\Omega$  and stipulate that  $\partial\Omega \cap \{\xi_j\}_{j=1}^K = \emptyset$ . In the above formulation (3),  $u_0(x)$  is the solution in the absence of heterogeneities and solves the calibration problem:

(4) 
$$\Delta u_0 = 0, \quad x \in \Omega; \qquad \partial_{\nu} u_0 = g, \quad x \in \partial \Omega; \qquad \int_{\partial \Omega} g = 0.$$

The condition  $\int_{\partial\Omega} g = 0$  is required for equation (4) to admit a consistent solution. The other key constituents of (3) are the polarisation matrices  $\{\mathcal{M}_j\}_{j=1}^K$  which are determined by the shape of each  $\bar{D}_j$  and N(x,y), the Neumann function for  $\Omega$ . The asymptotic formulation is derived with the assumption that the inclusions are well separated in the sense that  $|\xi_i - \xi_j| = \mathcal{O}(1)$  for  $i \neq j$  as  $\varepsilon \to 0$ . The problem of locating multiple closely spaced small inclusions was studied in [24] and an algorithm developed to detect some of their geometric features from an effective polarisation tensor.

As a consequence of (3), the following problem becomes central to our considerations;

**Problem 1:** Suppose that u(x) satisfies:

(5) 
$$\begin{cases} \Delta u = \sum_{j=1}^{K} w_j \cdot \nabla \delta(x - \xi_j), & x \in \Omega; \\ u(x) = f(x), & \partial_{\nu} u(x) = g(x), & x \in \partial \Omega. \end{cases}$$

Given f(x) and g(x), determine  $\xi_j$  and  $w_j$ ,  $j = 1 \dots K$ .

In fact, we describe a constructive algorithm to solve Problem 1 for three dimensional domains  $\Omega$ . The points  $\xi_j$  recover the locations of the inclusions whereas the vectors  $w_j$  are related to the shape of the inclusion via the polarisation matrix. Specialising to the case where  $D_j = B_{\varepsilon_j}(\xi_j)$  is a ball of a small radius  $\varepsilon_j$  centred at  $\xi_j$ , the explicit computation of the polarisation tensor yields [13],

(6) 
$$w_{j} = \begin{cases} 2\pi \nabla u_{0}(\xi_{j}) \frac{1 - \kappa_{j}}{1 + \kappa_{j}} \varepsilon_{j}^{2}, & \text{in } \mathbb{R}^{2}; \\ 4\pi \nabla u_{0}(\xi_{j}) \frac{1 - \kappa_{j}}{2 + \kappa_{j}} \varepsilon_{j}^{3}, & \text{in } \mathbb{R}^{3}. \end{cases}$$

This formula will be used in numerical tests in §3 below.

As was done for two dimensions in [9], in three dimensions, Problem 1 is reduced to finding K roots of a certain complex polynomial of degree K, plus solving several linear systems. In this sense, the algorithm is deterministic (non-iterative except for finding the roots of a K-th degree polynomial, which is a standard numerical procedure). Note that in three dimensions, there are 6K unknowns (three components for each of  $\xi_j$  and  $w_j$ ) that must be recovered. This makes a deterministic method especially valuable, since iterative methods often have convergence problems when faced with so many unknowns.

We remark that in [25], the authors considered a more general problem than Problem 1 in two dimensions; namely, they considered sources that consisted of both dipoles and monopoles. However, only dipoles arise in the asymptotic reduction of (1); so we do not include the monopoles in (5).

The efficacy of these algorithms is demonstrated on a number of sample problems comprised of multiple inclusions inside non-spherical domains. We employ the finite element solver FlexPDE to obtain a numerical solution to the forward problem (1) and the boundary data  $f(x) = u|_{\partial\Omega}$  for general regions. Numerical experiments confirm the relative robustness of our algorithm (typical values we tried are K=4 with 24 unknowns to solve for in  $\mathbb{R}^3$ ). Interestingly, despite the fact that the algorithm is designed to find small ball-like inclusions, it can provide information on the shape of more general small inclusions through approximation by multiple ball-like inclusions. We now state our main results.

1.1. **Description of the algorithm.** We now state the algorithm in three dimensions (In Appendix A, we include Algorithm 2, the solution to Problem 1 in the 2D case [9]).

Algorithm 1. (Solution to Problem 1 in  $\mathbb{R}^3$ ).

1. Compute the 2K complex integrals

(7) 
$$A_n = \int_{\partial\Omega} f(y)(y_1 + iy_2)^{n-1} (\nu_1 + i\nu_2) dS(y) - \int_{\partial\Omega} \frac{1}{n} g(y)(y_1 + iy_2)^n dS(y), \quad n = 1 \dots 2K$$

where  $\nu = (\nu_1, \nu_2, \nu_3)$  is the normal vector, and  $i = \sqrt{-1}$ .

2. Solve the system

(8) 
$$A_n = \sum_{j=1}^K \omega_j \eta_j^{n-1}, \quad n = 1 \dots 2K$$

using Lemma 1 (see below) to recover  $\omega_j$ ,  $\eta_j$ . The first two components of  $\xi_j$ ,  $w_j$  are then given by  $(\xi_{j1}, \xi_{j2}) = (\text{Re}(\eta_j), \text{Im}(\eta_j))$  and  $(w_{j1}, w_{j2}) = (\text{Re}(\omega_j), \text{Im}(\omega_j))$ .

3. Compute the 2K complex integrals

$$B_{n} = \int_{\partial\Omega} f(y) \left\{ (n-1)(y_{1} + iy_{2})^{n-2}y_{3}(\nu_{1} + i\nu_{2}) + (y_{1} + iy_{2})^{n-1}\nu_{3} \right\} dS(y)$$

$$- \int g(y)(y_{1} + iy_{2})^{n-1}y_{3}dS(y), \qquad n = 1 \dots 2K.$$

4. Solve the linear system of 2K variables

(10) 
$$B_n = \sum_{j=1}^K (n-1)\eta_j^{n-2}\omega_j\xi_{j3} + \eta_j^{n-1}w_{j3}, \qquad n = 1\dots 2K$$

for the third components  $\xi_{j3}, w_{j3}, j = 1 \dots K$ .

We make several observations. There are three numerical tasks involved: computing 4K complex surface integrals (7,9); computing the roots of a K-th degree polynomial; and solving several linear systems of O(K) variables. To compute the integrals, one must discretise the surface of the domain using N meshpoints with  $N \gg K$ . The only nonlinear step is solving for K roots of a K-degree polynomial, which is a standard numerical routine. The total running time is then of O(KN) as computing the integrals is the most time-consuming of the three tasks, assuming that  $K \ll N$ . The number K itself can also be algorithmically determined, using the techniques described in [11]. Finally, we remark that the solution to (10) typically results in  $\xi_{j3}$  and  $w_{j3}$  having small imaginary parts due to imperfections in measurement; for the practical implementation, the small imaginary parts are simply discarded.

Algorithm 1 relies upon solving the system (8) for the 2K unknowns  $\omega_j$ ,  $\eta_j$ , j = 1...K. The requirement to solve such systems has arisen frequently in many recent works, (cf. [26, 22, 19, 11, 9, 21, 25]). A common and effective method to solve this problem, stated below, reduces the system (8) to finding the roots of a K-th degree polynomial. For a derivation, see [26].

**Lemma 1.** Consider the system of 2K equations

(11) 
$$c_n = \sum_{j=1}^K \alpha_j \beta_j^n, \quad n = 0 \dots 2K - 1.$$

for the 2K unknowns  $\{\alpha_j\}_{j=1}^K$ ,  $\{\beta_j\}_{j=1}^K$ . Then  $\{\beta_j\}_{j=1}^K$  are the K roots of the polynomial

(12) 
$$\beta^K + l_1 \beta^{K-1} + \ldots + l_{K-1} \beta + l_K = 0$$

where the coefficients  $l_1 \dots l_K$  satisfy the linear system:

(13) 
$$\begin{bmatrix} c_0 & c_1 & \dots & c_{K-1} \\ c_1 & c_2 & \dots & c_K \\ \vdots & & & \vdots \\ c_{K-1} & c_2 & \dots & c_{2K-2} \end{bmatrix} \begin{bmatrix} l_K \\ l_{K-1} \\ \vdots \\ l_1 \end{bmatrix} = \begin{bmatrix} -c_K \\ -c_{K+1} \\ \vdots \\ -c_{2K-1} \end{bmatrix}.$$

Once  $\{\beta_j\}_{j=1}^K$  are computed,  $\{\alpha_j\}_{j=1}^K$  are obtained by solving the linear system given by the first K equations of (11).

#### 2. Derivation of Algorithm 1.

Algorithm 1 combines the ideas of [9, 19, 21]. Given any harmonic function v(y) which satisfies  $\Delta v = 0$  inside  $\Omega$ , multiply (5) by v and integrate by parts. This yields the identity

(14) 
$$\int_{\partial\Omega} f(y)\nu \cdot \nabla v(y)dS(y) - \int_{\partial\Omega} g(y)v(y)dS(y) = \sum_{j=1}^{K} w_j \cdot \nabla v(\xi_j).$$

Next we take  $v(y) = |y|^n Y_n^m(\hat{y})$ , where  $\hat{y} = y/|y|$  and  $Y_n^m$  is the spherical harmonic, to obtain

(15) 
$$\int_{\partial\Omega} f(y)\nu \cdot \nabla_y (|y|^n Y_n^m(\hat{y})) dS(y) - \int_{\partial\Omega} |y|^n Y_n^m(\hat{y}) g(y) dS(y) = \sum_{j=1}^K w_j \cdot \nabla_\xi \left( |\xi_j|^n Y_n^m(\hat{\xi}_j) \right).$$

In principle, the expression (15) yields an arbitrarily large number of identities relating the unknowns  $w_j$ ,  $\xi_j$  to the (known) integrals. Following [19, 21], we first consider the particular sequence of these identities corresponding to m = n. Note that

$$|y|^n Y_n^n(\hat{y}) = C(y_1 + iy_2)^n$$

for a known constant C independent of y and that

$$\nu \cdot \nabla_y (|y|^n Y_n^n(\hat{y})) = Cn(y_1 + iy_2)^{n-1} (\nu_1 + i\nu_2).$$

Then, applying this expression to (15) with m = n yields

(16) 
$$\int_{\partial\Omega} \left[ f(y)n(y_1 + iy_2)^{n-1}(\nu_1 + i\nu_2) - g(y)(y_1 + iy_2)^n \right] dS(y) = \sum_{j=1}^K n\eta_j^{n-1}\omega_j.$$

where we used the notation  $\eta_j = \xi_{1j} + i\xi_{2j}$ ;  $\omega_j = w_{1j} + iw_{2j}$ . This identity is precisely the sequence given by (7-8), which determines the first two coordinates of  $\xi_j$  and  $w_j$ . To determine the third coordinate, we choose m = n - 1 in (15). From this, we obtain

$$|y|^n Y_n^{n-1}(\hat{y}) = C(y_1 + iy_2)^{n-1} y_3$$

which in turn gives

$$\nu \cdot \nabla_y \left( |y|^n Y_n^{n-1}(\hat{y}) \right) = C \left\{ (n-1)(y_1 + iy_2)^{n-2} y_3(\nu_1 + i\nu_2) + (y_1 + iy_2)^{n-1} \nu_3 \right\}.$$

This yields the second sequence of identities,

$$(17) \int_{\partial\Omega} f(y) \left\{ (n-1)(y_1 + iy_2)^{n-2} y_3(\nu_1 + i\nu_2) + (y_1 + iy_2)^{n-1} \nu_3 \right\} - g(y)(y_1 + iy_2)^{n-1} y_3 dS(y)$$

$$= \sum_{j=1}^K (n-1) \eta_j^{n-2} \xi_{3j} \omega_j + \eta_j^{n-1} w_{j3}.$$

These identities correspond to those presented in (7-10). The derivation now complete.

#### 3. Numerical experiments

We now perform several experiments to test the robustness of Algorithms 1 and 2. For all experiments, we have used the finite element solver FlexPDE to obtain the solution of the forward problem (1). This package uses adaptive mesh generation and local and global error controls. The user-specified error tolerance was chosen to be  $10^{-5}$ ; we verified that changing this tolerance did not change the results. For all of the examples below, we have used  $\partial_{\nu}u = g(x) = x_2$  as the boundary condition with internal conductivity  $\kappa = 19$  in (2). Once the constants  $w_j$  and  $\xi_j$  were recovered using Algorithm 1 or 2 (cf. Appendix A for Algorithm 2), the connection formulae (6) were then used to recover the inclusion radius.

The integrals such as (7) were computed as follows. First, the output values generated by FlexPDE on the boundary of the surface were exported to Matlab. These points were then used to generate a surface mesh using Matlab's convhulin function. For a typical example such as experiment 1 below, this resulted in a total of about 1800 meshpoints. The resulting surface mesh was used to generate the normals and surface elements and compute the required integrals.

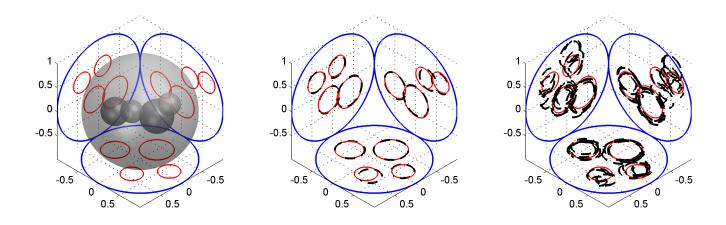


FIGURE 1. Recovery of four spherical inclusions inside a sphere (experiment 1). Left: the test configuration. Middle: no noise. Right: 2% noise.

**Experiment 1.** We take the domain to be unit ball with K = 4 spherical inclusions of varying size inside. The configuration and its recovery is shown in Figure 1; the sizes of four holes are

 $\varepsilon = 0.2,~0.2,~0.25,~0.3$ . Without any noise, excellent recovery is observed. To test the method's resilience to uncertainty in the measured data, we added 2% random perturbations to the output data of the forward solver. The right pane shows five distinct realisations of the recovery with noise. Note that the biggest inclusion exhibits the best recovery. Overall, very good agreement is observed.

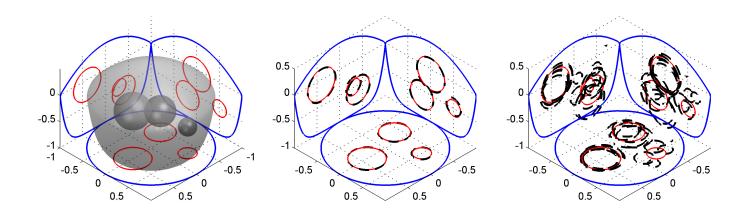


FIGURE 2. Recovery of three heterogeneities in a muffin-shaped domain (experiment 2). Middle: zero noise. Right: 2% noise.

**Experiment 2.** To demonstrate the efficacy of the recovery algorithm in more general regions, and not just spherical ones, we consider here a "muffin" shaped domain, as shown in Figure 2. It consists of the intersection of the top surface  $z = \frac{1}{2}(1 - (x^2 + y^2))$  and the bottom surface  $z = -(1 - (x^2 + y^2)^4)$ . Three spherical "raisins" of radii 0.15, 0.25, 0.3 are placed inside the "muffin". As in Experiment 1, the middle pane shows the recovery of the heterogeneity locations in the absence of noise, while the right pane shows five simulations of recovery with 2% random noise. Without noise, excellent recovery is observed. With noise, the biggest inclusion is still recovered relatively well.

**Experiment 3.** Consider a box-shaped anomaly inside a sphere as in Figure 3. We attempt its recovery by applying Algorithm 1 with K = 1, 2, 3. When K = 1, the recovered inclusion location captures well the centre of the box. Running the algorithm with K = 2, the two resulting inclusions approximate the box even better. Finally, K = 3 yields two inclusions which are identical to K = 2, plus a spurious inclusion of a very small radius (which appears as a smudge in the right pane). This indicates that the matrix in (13) is nearly singular when K = 3, and that K = 2 is a good place to stop.

**Experiments 4.** Figure 4 shows the recovery in two dimensions for a circular and a clover-shaped domain. The recovery is performed with no noise, 2% and 4% noise. For the noisy data, five different random realisations are shown. In the presence of noise, bigger inclusions are generally recovered better than the smaller ones. This is not surprising, since the smaller inclusions have less of an "imprint" on the boundary and become swamped as noise is increased.

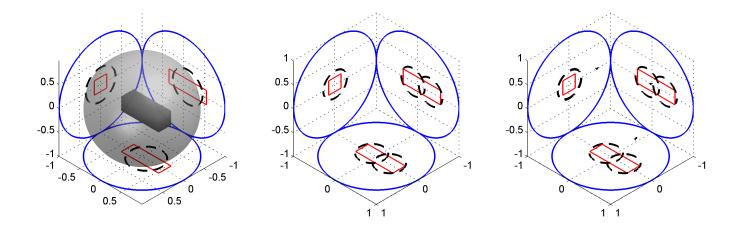


FIGURE 3. Recovery of a box-shaped inclusion inside a sphere using K = 1, 2, 3 disks (experiment 3). The original box is parallel to the axes and has corners at (0.15, 0.55, 0.15) and (-0.15, -0.35, -0.15).

**Experiment 5: L-shaped inclusion.** In the final experiment, we attempt to recover an L-shaped inclusion in Figure 5. The output of Algorithm 2 with  $K = 1 \dots 8$  is shown. Note that the corners are recovered surprisingly well when K = 8.

#### 4. Discussion

In this paper we tested an algorithm to recover K small inclusions inside a three-dimensional domain. It is based on the methods introduced in [9, 19, 21] for the two-dimensional case. It generalises the algorithm discussed in [23] in two ways: it is applicable for an arbitrary domains, and it is able to systematically handle multiple inclusions.

In [23], the authors write "given that these multi-source inverse problems are finite-dimensional, it may then be better to abandon the analytical approach and resort to a numerical method". We show that the numerical problem of identifying K dipoles in three dimensions for an arbitrary domain (which involves 6K unknowns), can be reduced to solving a single K degree polynomial, plus several systems of linear equations.

In [23], a recovery of a single source for Helmholtz equation was studied; the recovery of multiple sources was recently addressed for a single known wavelength was addressed in [27]. Unlike the EIT problem, the method in [27] only works in three dimensions and does not easily generalise to a two-dimensional setting. Recovering multiple sources for Helmholtz equation in two dimensions is still an open question.

In the presence of noise, higher spherical harmonics can easily get swamped. Since Algorithm 1 depends on spherical harmonics of degree l = 2K for the recovery of K dipoles, this makes it particularly prone to noise for larger values of K. In principle, there are l(l+1)/2 different spherical harmonics of degree at most l that can be used to recover the data; however Algorithm 1 only uses 2l of these. Therefore the algorithm is suboptimal as it does not use all of the available

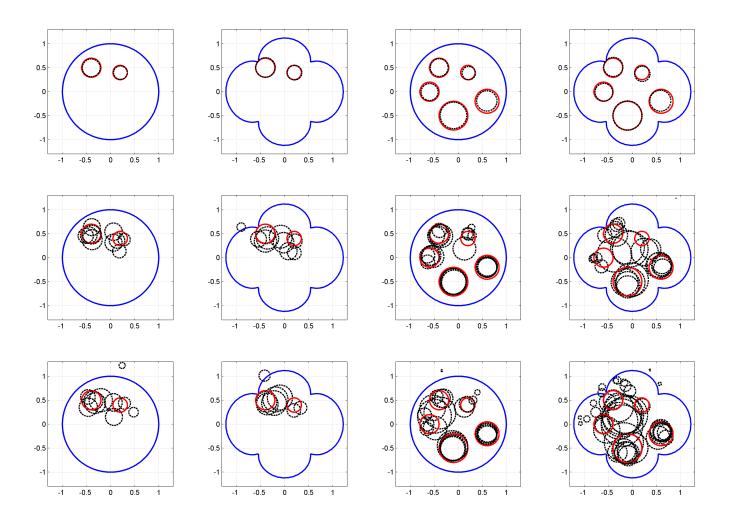


FIGURE 4. Recovery of two and five circular defects for a disk domain and for a "clover" domain (experiment 4). Dashed lines represent recovered defects using Algorithm 2 (cf. Appendix A). Solid lines (red online) give the original defect locations. First row: No noise. Second row: 2% noise in the output function f(x). Third row: 4% noise. Five different random simulations are shown for in each figure.

information. An optimal algorithm would only use harmonics of degree up to  $O(\sqrt{K})$  and would therefore be much less susceptible to noise. The design of such an optimal algorithm is an open problem. Another approach which can be explored to combat noisy data is oversampling: it is possible to improve recovery by generating more than the minimal number of equations. This leads to an overdetermined system; nonlinear error minimisation techniques could then be used to reduce the effect of the noise.

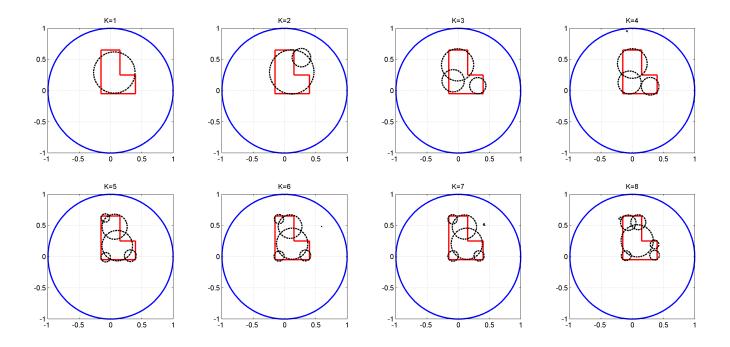


FIGURE 5. "Recovery" of an irregular L-shaped inclusion using K discs (experiment 5).

Numerical experiments with irregular-shaped inclusions suggest that this approach can be useful for recovery of larger contiguous patches (see Figure 5). These appear to be approximated surprisingly well by multiple inclusions, even though they are not separated. A related approach was used in [28] to recover inclusions consisting thin segments.

In our experiments, we fixed the number of inclusions K à priori and then applied the algorithm. As K is increased, the matrix (13) often becomes ill-conditioned, resulting in tiny inclusions (i.e. nearly zero radius). This is visible for example in figure 3 when K=3 and in figure 5 when K=4,6,7. Detecting when the matrix (13) becomes singular (by thresholding on determinant or condition number, for instance) is a good way to determine the appropriate K. This is the approach that was also used in [11] for example. Surprisingly, as figure 5 shows, the matrix (13) can be ill-conditioned for say K=4 but becomes well-conditioned again when K=5. This deserves further investigation.

#### 5. Acknowledgements

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## APPENDIX A: TWO DIMENSIONS

Here we recall the algorithm in two dimensions which was derived in [9].

# Algorithm 2. ([9], Solution to Problem 1 in $\mathbb{R}^2$ ).

1. Compute the 2K complex integrals

(18) 
$$A_n = \int_{\partial \Omega} f(y) \, y^{n-1} \, \nu \, dS(y) - \frac{1}{n} \int_{\partial \Omega} g(y) \, y^n \, dS(y), \quad n = 1 \dots 2K$$

where y and  $\nu$  are to be interpreted as complex numbers with  $y^{n-1}\nu$  corresponding to complex multiplication.

2. Solve the system

(19) 
$$A_n = \sum_{j=1}^K w_j \xi_j^{n-1}, \qquad n = 1 \dots 2K$$

using the Lemma 1 to recover  $w_i, \xi_i, j = 1 \dots K$ .

We now show that the method used in [12] for a single inclusion is equivalent to this algorithm; this also provides its alternative derivation. We summarise the method used in [12]. It is based on the function originally defined in [10],

(20) 
$$H(x) = \int_{\partial\Omega} f(y)\nu \cdot \nabla_y \Phi(x, y) dS(y) - \int_{\partial\Omega} g(y)\Phi(x, y) dS(y)$$

where

$$\Phi(x,y) = \frac{1}{2\pi} \ln|x - y|$$

is the free-space Green's function,  $\Delta \Phi = \delta(x-y)$ . It then follows from (14) that for  $x \notin \Omega$ ,

(21) 
$$H(x) = \sum_{j=1}^{K} w_j \cdot \nabla_{\xi} \Phi(x, \xi_j), \quad x \notin \Omega.$$

The idea in [12] is to take a constant R such that  $x \notin \Omega$  whenever |x| = R, and then look at the Fourier decomposition of  $H(Re^{i\theta})$ ; then solve the resulting system for  $\xi_j, w_j$ . We now show that this is equivalent to Algorithm 2 (cf. Appendix A). Indeed the following identity holds,

(22) 
$$\ln|x - y| = -\sum_{n=1}^{\infty} \frac{1}{2n} \left( x^{-n} y^n + \overline{x^{-n} y^n} \right) + \ln|x|, \qquad |x| > |y|$$

where  $x, y \in \mathbb{C}$  are complex numbers. Take  $x = Re^{i\theta}$  and compute using (22),

(23) 
$$\int_0^{2\pi} \Phi(Re^{i\theta}, y)e^{in\theta}d\theta = -\frac{1}{4n}R^{-n}y^n.$$

Next, note that for any vector w we have

$$w \cdot \nabla y^n = n (y_1 + iy_2)^{n-1} (\hat{\imath} + i\hat{\jmath}) \cdot w$$
$$= n y^{n-1} w$$

so that, in combination with (22)

(24) 
$$\int_0^{2\pi} w \cdot \nabla_y \Phi(Re^{i\theta}, y) e^{in\theta} d\theta = -\frac{1}{4} R^{-n} y^{n-1} w.$$

Next, choose R sufficiently large such that  $x \notin \Omega$  for whenever |x| = R. Integrate both sides of (14) with respect to x on a circle of radius R. Using formulae (23) and (24), we then obtain

$$\int_{\partial\Omega} n f(y) y^{n-1} \nu dS(y) - \int_{\partial\Omega} g(y) y^n dS(y) = \sum_{j=1}^K n \xi_j^{n-1} w_j$$

This is corresponds precisely to (18, 19). The derivation is complete.

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