

Homework 5: Non-context-free languages

CSE 30151 Spring 2018

Due 2018/03/22 at 10:00pm

Instructions

- Create a PDF file (or files) containing your solutions. You can write your solutions by hand, but please scan them in the library or using a smartphone to get them into PDF.
- Please name your PDF file(s) as follows to ensure that the graders give you credit for all of your work:
 - If you're making a complete submission, name it *netid-hw5.pdf*, where *netid* is replaced with your NetID.
 - If you're submitting some problems now and want to submit other problems later, name it *netid-hw5-123.pdf*, where 123 is replaced with the problem numbers you are submitting at this time.
- Submit your PDF file(s) in Sakai. Don't forget to click the Submit button!

Problems

Each problem is worth 10 points.

1. **Pumping lemma for CFLs.** Choose **two out of three** of the following languages and use the pumping lemma to show that they are not context free. If you do all three, we'll give credit for the best two.

(a) [Problem 2.30a] $A = \{0^n 1^n 0^n 1^n \mid n \geq 0\}$.

(b) [Problem 2.31]

$B = \{w \in \{0, 1\}^* \mid w = w^R \text{ and } w \text{ has an equal number of 0s and 1s}\}$.

(c) [cf. HW3, Q2b] The language over $\Sigma = \{0, 1, +, =\}$:

$ADD = \{x = y + z \mid x, y, z \text{ are binary natural numbers, and } x = y + z\}$.

2. **The SCRAMBLE operation** [Problem 2.43]. If w and w' are strings over an alphabet Σ , define the relation $w \doteq w'$ to be true iff w' is a permutation of w , that is, they have the same number of each type of symbol, but possibly in a different order. If w is a string and L is a language, define

$$\begin{aligned}\text{SCRAMBLE}(w) &= \{w' \mid w' \doteq w\} \\ \text{SCRAMBLE}(L) &= \bigcup_{w \in L} \text{SCRAMBLE}(w).\end{aligned}$$

For example,

$$\begin{aligned}\text{SCRAMBLE}(0^*1^*) &= \{0, 1\}^* \\ \text{SCRAMBLE}((01)^*) &= \{w \in \{0, 1\}^* \mid w \text{ has an equal number of 0s and 1s}\}.\end{aligned}$$

- (a) Let $\Sigma = \{\mathbf{a}, \mathbf{b}\}$. Show that for any regular language L over Σ , $\text{SCRAMBLE}(L)$ is context-free.
- (b) Let $\Sigma = \{\mathbf{a}, \mathbf{b}, \mathbf{c}\}$. Show that there exists a regular language over Σ such that $\text{SCRAMBLE}(L)$ is *not* context-free.

3. Non-closure properties

- (a) [Exercise 2.2a] Use the languages

$$\begin{aligned}A &= \{\mathbf{a}^m \mathbf{b}^n \mathbf{c}^n \mid m, n \geq 0\} \\ B &= \{\mathbf{a}^n \mathbf{b}^n \mathbf{c}^m \mid m, n \geq 0\}\end{aligned}$$

to prove that context-free languages are *not* closed under intersection.

- (b) [Exercise 2.2b] Use (a) and DeMorgan's law to prove that context-free languages are *not* closed under complementation.