## Homework 2: DFAs and NFAs

Theory of Computing (CSE 30151), Spring 2024

## Due 2024-02-02 5pm

## Instructions

- Create a PDF file (or files) containing your solutions. You can write your solutions by hand, but please scan them into a PDF.
- Please name your PDF file(s) as follows to ensure that the graders give you credit for all of your work:
  - If you're making a complete submission, name it netid-hw2.pdf, where netid is replaced with your NetID.
  - If you're submitting some problems now and want to submit other problems later, name it *netid*-hw2-part123.pdf, where 123 is replaced with the problem numbers you are submitting at this time.
- Submit your PDF file(s) in Canvas.

## Problems (10 points each)

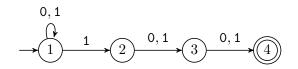
1. Divisibility tests. Define, for all k > 0,

 $D_k = \{w \in \{0, \dots, 9\}^* \mid w \text{ is the decimal representation of a multiple of } k\}$ 

where  $\varepsilon$  is considered to represent the number 0. For example, the strings  $\varepsilon$ , 0, 88, and 088 all belong to  $D_2$ , but 99 and 099 do not.

- (a) Prove that  $D_2$  is regular by writing a DFA for  $D_2$ .
- (b) Prove that  $D_3$  is regular by writing a DFA for  $D_3$ .
- (c) Prove that  $D_6$  is regular. An explicit DFA is not necessary.
- (\*) Optional alternative: You can get full credit for all of the above if you can prove that for any k > 0,  $D_k$  is regular, by describing how to write the formal description of a DFA  $M = (Q, \{0, \ldots, 9\}, \delta, s, F)$  in terms of k.

2. Nondeterminism. Consider the following NFA  $N_2$  (same as in Figure 1.31), which accepts a string iff the third-to-last symbol is a 1:



- (a) Use the subset construction (Theorem 1.39) to convert  $N_2$  to a DFA M. You may omit curly braces and commas when naming states; for example, instead of  $\{1, 2, 3, 4\}$  you may write 1234. (Hint: the DFA should be equivalent to the one in Figure 1.32.)
- (b) Why are the states in Figure 1.32 named  $q_{abc}$  where  $a, b, c \in \{0, 1\}$ ?
- (c) In Example 1.30, Sipser asks what happens if you modify  $N_2$  into the following NFA let's call it  $N'_2$ :

$$\xrightarrow{0,1} \xrightarrow{1} 2 \xrightarrow{0,1,\varepsilon} 3 \xrightarrow{0,1,\varepsilon} 4$$

Describe in English what language  $N'_2$  recognizes.

- (d) Use the subset construction (Theorem 1.39) to convert  $N'_2$  to a DFA M'.
- 3. Procrustean closure properties. Let  $\Sigma$  be an alphabet, and let  $L_3 = \{\text{theory}, \text{of}, \text{computing}\}$  be an example language.
  - (a) For any  $w = w_1 w_2 \cdots w_{n-1} w_n$ , define

 $STRETCH(w_1w_2\cdots w_n) = w_1w_1w_2w_2\cdots w_{n-1}w_nw_n.$ 

This induces an operation on languages,

 $STRETCH(L) = {STRETCH(w) | w \in L}.$ 

For example,

 $STRETCH(L_3) = \{thheeorryy, ooff, ccoommpputtiinngg\}.$ 

Prove that if L is a regular language, then STRETCH(L) is also regular.

(b) For any  $w = w_1 w_2 \cdots w_{n-1} w_n$  with  $n \ge 2$ , define

 $CHOP(w_1w_2\cdots w_{n-1}w_n)=w_2\cdots w_{n-1}.$ 

This induces an operation on languages,

 $CHOP = \{CHOP(w) \mid w \in L \text{ and } |w| \ge 2\}.$ 

For example,

$$CHOP(L_3) = \{\texttt{heor}, \varepsilon, \texttt{omputin}\}.$$

Prove that if L is a regular language, then CHOP(L) is also regular.