

Homework 1: Strings and languages

CSE 30151 Fall 2020

Due: 2020/08/21 5:00pm

Instructions

- Create a PDF file (or files) containing your solutions. You can write your solutions by hand, but please scan them into a PDF.
- Please name your PDF file(s) as follows to ensure that the graders give you credit for all of your work:
 - If you're making a complete submission, name it *netid-hw1.pdf*, where *netid* is replaced with your NetID.
 - If you're submitting some problems now and want to submit other problems later, name it *netid-hw1-123.pdf*, where 123 is replaced with the problem numbers you are submitting at this time.
- Submit your PDF file(s) in Sakai. Don't forget to click the Submit button!

Problems (10 points each)

1. Proof practice.

- (a) Convert this paragraph proof to a statement-reason proof. Please be sure to write which statement(s) each statement depends on.

To show: If s is a string, every substring of a substring of s is a substring of s .

Proof: Let y be a substring of s , that is, $s = xyz$ for some x, z ; and let v be a substring of y , that is, $y = uvw$ for some u, w . Then $s = xuvwz$, so v is a substring of s .

- (b) Convert this statement-reason proof to a paragraph proof.

To show: If w is a string, every prefix of a suffix of w is a suffix of a prefix of w .

1. v is a suffix of w Given
2. y is a prefix of v Given
3. $\exists x$ s.t. $xv = w$ (1), def. suffix
4. $\exists z$ s.t. $yz = v$ (2), def. prefix
5. $xyz = w$ (3), (4), substitution
6. xy is a prefix of w (5), def. prefix
7. y is a suffix of xy (6), def. suffix

2. **String homomorphisms.** If Σ and Γ are finite alphabets, define a *string homomorphism* to be a function $\phi : \Sigma^* \rightarrow \Gamma^*$ that has the property that for any $u, v \in \Sigma^*$, $\phi(uv) = \phi(u)\phi(v)$.

For example, the function $\phi : \{0, \dots, 9, \mathbf{A}, \dots, \mathbf{F}\}^* \rightarrow \{0, 1\}^*$ that converts a hexadecimal number with $n \geq 0$ digits into a binary number with $4n$ bits is a string homomorphism:

$$\begin{aligned}\phi(\varepsilon) &= \varepsilon \\ \phi(0) &= 0000 \\ \phi(\mathbf{A}) &= 1010 \\ \phi(\mathbf{CAB}) &= 110010101011\end{aligned}$$

Intuitively, a string homomorphism does a “search and replace” where each symbol is replaced with a (possibly empty) string. Prove this more formally: that is, prove that if ϕ is a string homomorphism, then for any $w = w_1 \cdots w_n$ (where $n \geq 0$ and $w_j \in \Sigma$ for $1 \leq j \leq n$), we have

$$\phi(w) = \phi(w_1) \cdots \phi(w_n). \quad (*)$$

Use induction on n .

- (a) State and prove the base case ($n = 0$).
 - (b) Assume that (*) is true for $n = i$ and prove (*) for $n = i + 1$.
3. **Finite and cofinite.** For any alphabet Σ , let FINITE be the set of all finite languages over Σ , and let

$$\text{coFINITE} = \{L \mid \bar{L} \in \text{FINITE}\}$$

(where, for any language L over Σ , \bar{L} is the complement of L , that is, $\Sigma^* \setminus L$). For example, Σ^* is in coFINITE because its complement is \emptyset , which is finite. (Please think carefully about this definition, and note that coFINITE isn't the same thing as $\overline{\text{FINITE}}$.)

- (a) Are there any languages over Σ in $\text{FINITE} \cap \text{coFINITE}$? Prove your answer.
- (b) Are there any languages over Σ that are *not* in $\text{FINITE} \cup \text{coFINITE}$? Prove your answer.