Linear programming, graphically

We've seen examples of problems that lead to linear *constraints* on some unknown quantities.

Now we are going to add an extra ingredient: some quantity that we want to maximize or minimize, such as profit, or costs.

If the quantity to be maximized/minimized can be written as a linear combination of the variables, it is called a linear *objective function*.

Linear programming is the business of finding a point in the feasible set for the constraints, which gives an optimum value (maximum or a minimum) for the objective function.

We'll see how a linear programming problem can be solved graphically.

Example — constraints

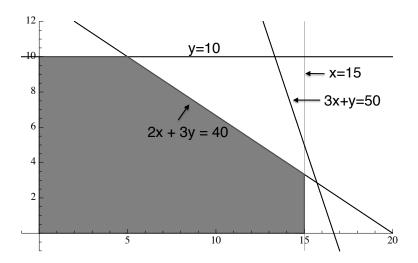
A juice stand sells two types of fresh juice in 12 oz cups, the Refresher and the Super-Duper. The Refresher is made from 3 oranges, 2 apples and a slice of ginger. The Super Duper is made from one slice of watermelon, 3 apples and one orange. The owners of the juice stand have 50 oranges, 40 apples, 10 slices of watermelon and 15 slices of ginger. Let x denote the number of Refreshers they make and let y denote the number of Super-Dupers they make.

Last time, we saw that the set of constraints on x and y was:

$$3x + y \le 50$$
 $2x + 3y \le 40$
 $x \le 15$ $y \le 10$
 $x \ge 0$ $y \ge 0$

Example — the feasible set

Here is the *feasible set*, the set of combinations of x and y that are possible given the limited supply of ingredients:



Example — adding an objective

Now suppose that Refreshers sell for \$6 each and Super-Dupers sell for \$8 each. Let's suppose also that the juice stand will sell all of the drinks they can make on this day, so their revenue for the day is 6x + 8y. If a goal of the juice stand is to maximize revenue, then they want to maximize the value of 6x + 8y, given the constraints on production.

In other words they want to find a point (x,y) in the feasible set which gives a maximum value for the **objective** function 6x + 8y. [Note that the value of the objective function (6x + 8y = revenue) varies as (x,y) varies over the points in the feasible set. For example if (x,y) = (2,5), revenue = 6(2) + 8(5) = \$52, whereas if (x,y) = (5,10), revenue = 6(5) + 8(10) = \$110.]

Terminology

Suppose we are given a problem that involves assigning values $x,\,y$ to some quantities.

The choices of x, y may be subject to some *constraints*: linear inequalities of the form

$$a_0x + a_1y \le b,$$
 $a_0x + a_1y < b,$ $a_0x + a_1y \ge b,$ $a_0x + a_1y > b,$

where a_0 , a_1 and b are constants.

There is a linear objective function: an expression of the form cx + dy, where c and d are constants, and we wish to find the maximum or minimum value that the objective function can take on the feasible set. We use the term *optimal value* to cover both maximizing and minimizing.

A linear programming problem is the problem of finding a point $(x_0, y_0) \in F$, the feasible set where all constraints are satisfied, with $O(x_0, y_0)$ as big as possible (if we are doing a maximum problem), or as small as possible (if we are minimizing).

A BIG IDEA of linear programming

If the feasible set of a linear programming problem is bounded (contained inside some big circle; equivalently, there is no direction in which you can travel indefinitely while staying in the feasible set), then, whether the problem is a minimization or a maximization, there will be an optimum value. Furthermore:

- there will be some *corner point* of the feasible region that is an optimum
- if there is more than one optimum corner point then there will be exactly two of them, they will be adjacent, and any point in the line between them will also be optimum.

The picture on the next page, taken from page 238 of the text, illustrates this graphically.

A BIG IDEA of linear programming

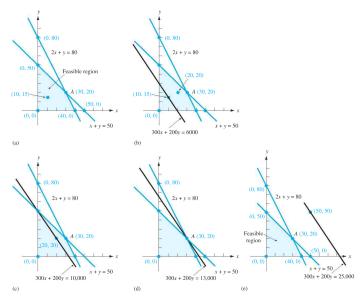
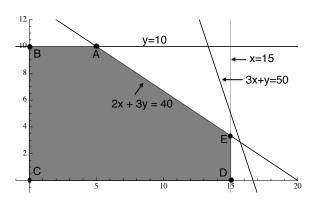


FIGURE 3-16 As an objective function increases, the objective function line moves toward a corner. In (e) the objective function moves out of the feasible region and out of consideration.

We now know that the maximum of 6x + 8y on the feasible set occurs at a corner of the feasible set since the feasible set is bounded (it may occur at more than one corner, but it occurs at at least one). We already have a picture of the feasible set and below, we have labelled the corners, A, B, C, D and E.



To find the maximum value of 6x + 8y and a point (x, y) in the feasible set at which it is achieved, we need only calculate the co-ordinates of the points A, B, C, D and E and compare the value of 6x + 8y at each.

Point	Coordinates	Value of $6x + 8y$
A		
В		
С	(0, 0)	0
D		
Е		

Point	Coordinates	Value of $6x + 8y$
A	(5,10)	110
В	(0,10)	80
C	(0, 0)	0
D	(15,0)	90
Е	(15,10/3)	$\frac{350}{3}$

Hence E is the largest value (and C is the smallest).

The way to generate optimal revenue is to produce 15 Refreshers and 3.33... Super-Dupers, for a total revenue of \$116.66

BUT: this solution is **not** feasible, since we can only produce whole-number quantities of drinks. We might round down y (to make sure we stay in the feasible region), and report an answer of 15 Refreshers and 3 Super-Dupers, for a total revenue of \$114. If we explore a little more, we find that x = 14 and y = 4 is also feasible, and yields revenue \$116. This is the optimum.

In general, dealing with integral constraints is **very** tough. In practice, if there are integral constraints, the best way to proceed is to solve the problem without worrying about this additional condition and then look at the nearby integer points to find the maximum.

Example with many solutions

Suppose that Super-Dupers sell for \$9 instead of \$8. What is the new maximum revenue?

Now our revenue function is 6x + 9y. The feasible set remains unchanged, so the corners remain unchanged. We test the new objective at each corner, as before:

Point	Coordinates	Value of $6x + 9y$
A	(5,10)	120
В	(0,10)	90
C	(0, 0)	0
D	(15,0)	90
E	(15,10/3)	120

Now there are two

corners — A and E — that have maximum revenue, \$120. Both of these are solutions, as is any point that lies in a straight line between them.

There is a whole line of solutions because the new revenue line 6x + 9y is parallel to the constraint line 2x + 3y = 40, that joins A and E.

If the feasible set is not bounded

If the feasible set of a linear programming problem is not bounded (there is a direction in which you can travel indefinitely while staying in the feasible set) then a particular objective may or may not have an optimum:

- if it is a maximization problem, there might be a maximum, or it might be possible to make the objective arbitrarily large inside the feasible set, and
- if it is a minimization problem, there might be a minimum, or it might be possible to make the objective arbitrarily small (big and negative) inside the feasible set.

If the feasible set is not bounded

If there is a maximum/minimum, it can happen

- uniquely at a corner,
- at two adjacent corners and at all points in between,
- at a corner or along an infinite ray leaving that corner, or
- along an entire infinite line.

This is a little more nuanced than the Theorem stated on page 239 of the text (which is not really a true theorem ☺).

The next slide present some examples.

Some unusual examples

Suppose the only constraint is $y \ge 0$. Then the feasible set is unbounded and has no corners.

- Objective = y has no maximum. It has a minimum, reached along the entire x-axis.
- Objective = -y has no minimum, but has a maximum
- ▶ Objective = x y has no minimum, and no maximum

Suppose the constraints are $y \ge 0$ and $x \ge 0$. Then the feasible set is unbounded and has one corner.

- Objective = x + y has a minimum, reached uniquely at the corner.
- Objective = y has a minimum, reached along the ray starting at the corner and moving to the right.

Suppose the constraints are $y \ge 0$, $x \ge 0$, $y \le 2$. Then the feasible set is unbounded and has two corners.

• Objective = x has a minimum, reached at both corners, and between the two corners.

A diet example

Mr. Carter eats a mix of Cereal A and Cereal B for breakfast. The amount of calories, sodium and protein per ounce for each is shown in the table below. Mr. Carter's breakfast should provide at least 480 calories but less than or equal to 700 milligrams of sodium. Mr. Carter would like to maximize the amount of protein in his breakfast mix.

	Cereal A	Cereal B
Calories (per ounce)	100	140
Sodium(mg per ounce)	150	190
$\mathbf{Protein}(\mathbf{g}\ \mathbf{per}\ \mathbf{ounce})$	9	10

Let x denote the number of ounces of Cereal A that Mr. Carter has for breakfast and let y denote the number of ounces of Cereal B that Mr. Carter has for breakfast. We find that the set of constraints for x and y are

 $100x + 140y \ge 480$, $150x + 190y \le 700$, $x \ge 0$, $y \ge 0$

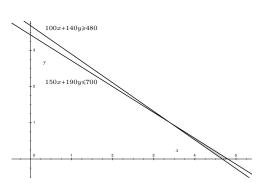
A diet example

(a) What is the objective function?

Objective = 9x + 10y.

(b) Graph the feasible set.

$$100x + 140y \ge 480$$
, $150x + 190y \le 700$, $x \ge 0$, $y \ge 0$



The feasible set is the skinny triangle just above the x-axis.

A diet example

(c) Find the corners of the feasible set and the maximum of the objective function on the feasible set.

The two corners on the x-axis are (4.8,0) and $(\frac{16}{3},0)$. The intersection of 100x + 140y = 480 and 150x + 190y = 700 is the point $(\frac{17}{5},1)$. The values of the objective function are 42.2, 48 and $\frac{203}{5} = 40.6$. Hence the maximum of the objective function is 48 and it occurs at the point $(\frac{16}{3},0)$ on the boundary of the feasible region and nowhere else. (The minimum occurs at 40.6 where the two lines intersect).

(d) Conclusion?

To maximize intake of protein while keeping sodium and calories at acceptable levels, Mr. Carter should eat 5.33... ounces of Cereal A and none of cereal B.

Michael is taking an exam in order to become a volunteer firefighter. The exam has 10 essay questions and 50 multiple choice questions. He has 90 minutes to take the exam. The essay questions are worth 20 points each and the short-answer questions are worth 5 points each. An essay question takes 10 minutes to answer and a short-answer question takes 2 minutes. Michael must do at least 3 essay questions and at least 10 short-answer questions. Michael knows the material well enough to get full points on all questions he attempts and wants to maximize the number of points he will get.

Let x denote the number of multiple choice questions that Michael will attempt and let y denote the number of essay questions that Michael will attempt. Write down the constraints and objective function in terms of x and y and find the/a combination of x and y which will allow Michael to gain the maximum number of points possible.

```
2x + 10y \le 90 (time needed to answer the questions).

x \ge 10 (at least 10 short-answer questions).

x \le 50 (at most 50 short-answer questions).

y \ge 3 (at least 3 essay questions).

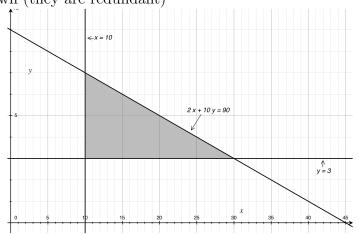
y \le 10 (at least 3 essay questions).

5x + 20y is the objective function (Michael's total score).

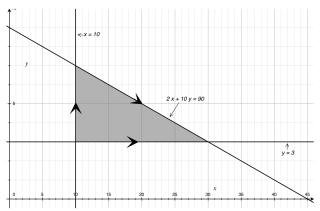
It is required to maximize the objective.
```

Here is the feasible region.

The lines x = 50 and y = 10 are outside the part of the plane shown (they are redundant)



The three corners are (10,3), (10,7) and (30,3). The values of the objective function at these corners are 60, 190 and 210. Hence Michael can maximize his score by answering 3 essay questions and 30 short-answer questions.



Example with unbounded region

A local politician is budgeting for her media campaign. She will distribute her funds between TV ads and radio ads. She has been given the following advice by her campaign advisers;

- She should run at least 120 TV ads and at least 30 radio ads.
- The number of TV ads she runs should be at least twice the number of radio ads she runs but not more than three times the number of radio ads she runs.

The cost of a TV ad is \$8000 and the cost of a radio ad is \$2000. Which combination of TV and radio ads should she choose to minimize the cost of her media campaign?

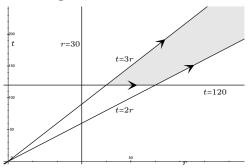
Let t be the number of TV ads and r the number of radio ads.

$$t \geqslant 120$$
 $r \geqslant 30$ $2r \leqslant t \leqslant 3r$

The objective function is 8000t + 2000r which she wishes to minimize.

Example with unbounded region

Here is the feasible region, which we see is unbounded.



If we wanted to maximize the campaign cost, we have no trouble; buy r radio ads and anywhere between 2r and 3r television ads, for a cost of between \$6,000r and \$8,000r — and r can be as large as we like.

The minimum cost occurs at the vertex where t = 120 and t = 3r meet. This is the point r = 40, t = 120. Hence she should buy 40 radio ads and 120 TV ads.

Example with no feasible region

Mr. Baker eats a mix of Cereal A and Cereal B for breakfast. The amount of calories, sodium and protein per ounce for each is shown in the table below. Mr. Baker's breakfast should provide at least 600 calories but less than 700 milligrams of sodium. Mr. Baker would like to maximize the amount of protein in his breakfast mix.

	Cereal A	Cereal B
Calories (per ounce)	100	140
Sodium(mg per ounce)	150	190
Protein(g per ounce)	9	10

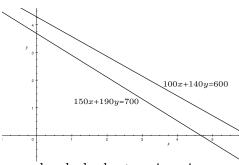
I am I am I

Let x denote the number of ounces of Cereal A that Mr. Baker has for breakfast and let y denote the number of ounces of Cereal B that Mr. Baker has for breakfast.

Example with no feasible region

(a) Show that the feasible set is empty here, so that there is no feasible combination of the cereals for Mr. Baker's breakfast.

$$100x + 140y \ge 600$$
 (calories)
 $150x + 190y < 700$ (sodium)
 $x \ge 0, y \ge 0$



Because we have $x \ge 0, \ y \ge 0$ we need only look at regions in the first quadrant.

Since (0,0) satisfies 150x + 190y < 700, $P_0 = \{1\}$. Since (0,0) satisfies 100x + 140y < 600, $P_1 = \emptyset$.

Example with no feasible region

(b) If Mr. Baker goes shopping for new cereals, what should he look for on the chart giving the nutritional value, so that he can have some feasible combination of the cereals for breakfast?

This is an essay question with no single right answer. Basically Mr. Baker needs to choose cereals with either more calories or less sodium per ounce.