

A PRIMAL-DUAL FORMULATION FOR CERTIFIABLE COMPUTATIONS IN SCHUBERT CALCULUS

JONATHAN D. HAUENSTEIN, NICKOLAS HEIN, FRANK SOTTILE

ABSTRACT. Formulating a Schubert problem as the solutions to a system of equations in either Plücker space or in the local coordinates of a Schubert cell typically involves more equations than variables. We present a novel primal-dual formulation of any Schubert problem on a Grassmannian or flag manifold as a system of bilinear equations with the same number of equations as variables. This formulation enables numerical computations in the Schubert calculus to be certified using algorithms based on Smale's α -theory.

INTRODUCTION

Numerical algebraic geometry provides fast, efficient methods to approximate all solutions to a system of polynomial equations [29]. When the system is *square* in that the number of equations is equal to the number of variables, Smale's α -theory [5, Ch. 8] gives methods to certify that these approximate solutions correspond to actual solutions. This is implemented in software which can be used to prove that all solutions have been found and to certifiably count the number of real solutions [12].

The Schubert calculus of enumerative geometry has come to mean all problems of determining the linear subspaces of a vector space that have specified positions with respect to other, fixed but general subspaces. Originating in work of Schubert [23, 24, 25, 26], it has long been an active subject, undergoing significant recent development. As a rich and well-understood class of geometric problems, the Schubert calculus is a laboratory for the systematic study of new phenomena in enumerative geometry [18, 32], particularly as Schubert problems are readily formulated and studied on a computer.

There are two traditional formulations of Schubert problems, one in global Plücker coordinates and one in the local coordinates of a Schubert cell [8]. The first involves quadratic Plücker equations and linear equations, while the second uses no Plücker equations and the linear equations are replaced by minors of matrices. Both formulations typically involve far more equations than variables. The second method was used to provide evidence for the Shapiro conjecture and discover its generalizations [7, 20, 21, 19, 31] through the systematic study of many billions of instances of several thousand Schubert problems [9, 11, 14, 22, 30]. These were exact computations, using symbolic methods. Numerical computation should enable the study of far more and far larger Schubert problems as well as more delicate monodromy computations on a computer. For this hope

2010 *Mathematics Subject Classification.* 14N15, 14Q20.

Key words and phrases. Schubert calculus, square systems, certification.

Research of Hauenstein supported in part by NSF grant DMS-1262428 and DARPA YFA. Research of Hein and Sottile supported in part by NSF grant DMS-0915211.

to become a reality, algorithms tailored to Schubert calculus need to be developed and implemented—this is starting to occur [15, 33]—and methods to certify computation need to be developed.

This second point was highlighted in a proof-of-concept paper to compute monodromy (Galois groups) of some Schubert problems [17]. This included numerically solving a Schubert problem with 17589 solutions and numerically computing enough monodromy permutations to conclude that the monodromy was the full symmetric group S_{17589} . These computations were not proofs in the ordinary sense as they relied upon software heuristics for their conclusions. This nevertheless inspired the development and implementation of methods to certify computations in numerical algebraic geometry.

Certification in numerical algebraic geometry rests upon work of Smale, who studied the convergence of iterations of Newton’s method applied to a point x_0 , considered to be an approximate solution to a system $F(x) = 0$ of polynomial equations, when the system is square [28]. This is called *α -theory* for the existence of an absolute constant $\alpha_0 > 0$ such that when a constant $\alpha = \alpha(x_0, F) > 0$ satisfies $\alpha < \alpha_0$, Newton iterates starting at x_0 converge quadratically (doubling significant digits with each step) to a solution to $F(x) = 0$. This was implemented in software [12] as an *a posteriori* test to certify the output of a numerical solver while Beltrán and Leykin [3, 4] developed and implemented a certified path-tracking algorithm which also certifies monodromy computations. In an important special case, *a posteriori* estimates suffice to certify monodromy [10].

Traditional formulations of Schubert problems typically lead to overdetermined systems (more equations than variables) so algorithms based on α -theory cannot be used to certify traditional numerical computation in Schubert calculus. We present a novel formulation of any Schubert problem in primal-dual coordinates as a square system of bilinear equations. This will enable certification based on α -theory, both *a posteriori* certification of approximate solutions and certified path-tracking.

The set of all ℓ -dimensional linear subspaces (ℓ -planes) in a vector space V having specified position with respect to a fixed subspace forms a *Schubert subvariety* of the Grassmannian. A Schubert problem is formulated as the intersection of a collection of Schubert varieties in a Grassmannian which are in general position. A natural extension is to consider a Schubert problem to be given by intersecting a collection of Schubert subvarieties in general position in a flag manifold. We give a square primal-dual formulation of any Schubert problem on any flag manifold.

The main ideas are well illustrated for the Grassmannian. In Section 1 we describe Schubert varieties in the Grassmannian and give a traditional formulation of a Schubert problem as an overdetermined system of determinantal equations in local coordinates for the Grassmannian. We introduce our primal-dual reformulation in Section 2, where we use duality to recast Schubert problems as a square system of bilinear equations in a larger space. We improve this, using a hybrid approach and more sophisticated local coordinates to our primal-dual formulation in Section 3. Finally, in Section 4 we explain how this extends to all Schubert problems in a flag manifold.

1. MODELING SCHUBERT PROBLEMS

We describe Schubert varieties, Schubert problems, and how to model them as systems of equations. For a positive integer n , write $[n]$ for the set $\{1, \dots, n\}$. Fix ℓ to be a positive integer at most $n-1$. Let $w_0 \in S_n$ be the permutation of n such that $w_0(i) = n+1-i$.

1.1. Schubert problems. Let V be an n -dimensional complex vector space. The *Grassmannian* $\text{Gr}(\ell, V)$ of ℓ -planes of V is a complex manifold of dimension $\ell(n-\ell)$ that is a subvariety of Plücker space, $\mathbb{P}(\wedge^\ell V) \simeq \mathbb{P}^{\binom{n}{\ell}-1}$.

The Grassmannian has distinguished Schubert subvarieties. A *flag* F_\bullet is a sequence

$$F_\bullet : \{0\} \subsetneq F_1 \subsetneq F_2 \subsetneq \dots \subsetneq F_n = V,$$

of linear subspaces of V where $\dim F_i = i$. The position of a ℓ -plane H with respect to the flag F_\bullet is the increasing sequence $\alpha : 1 \leq \alpha_1 < \dots < \alpha_\ell \leq n$ where

$$\alpha_i := \min\{j : \dim(H \cap F_j) = i\}, \text{ for } i = 1, \dots, \ell.$$

Write $\alpha := \alpha(H, F_\bullet)$ and call α a *Schubert condition*. All such increasing sequences may occur. Write $\binom{[n]}{\ell}$ for the set of Schubert conditions which is partially ordered by coordinatewise comparison: $\alpha \leq \beta$ if and only if $\alpha_i \leq \beta_i$ for each $i = 1, \dots, \ell$.

Let $X_\alpha^\circ F_\bullet$ be the set of all ℓ -planes H with $\alpha(H, F_\bullet) = \alpha$, which is a *Schubert cell*. This has dimension $\dim(\alpha) := (\alpha_1 - 1) + \dots + (\alpha_\ell - \ell)$. We have the Bruhat decomposition of the Grassmannian

$$\text{Gr}(\ell, n) = \bigsqcup_{\alpha \in \binom{[n]}{\ell}} X_\alpha^\circ F_\bullet.$$

Given a Schubert condition α and a flag F_\bullet , define the *Schubert subvariety* $X_\alpha F_\bullet$ by

$$(1.1) \quad X_\alpha F_\bullet := \{H \in \text{Gr}(\ell, V) \mid \dim(H \cap F_{\alpha_i}) \geq i, \text{ for } i = 1, \dots, \ell\},$$

which is the closure of the Schubert cell $X_\alpha^\circ F_\bullet$. This has dimension $\dim(\alpha)$ and thus codimension $|\alpha| := \ell(n-\ell) - \dim(\alpha) = \ell(n-\ell) - \sum_i (\alpha_i - i)$.

Example 1.1. Suppose that $\ell = 2$, $n = 8$, and F_\bullet is a complete flag in \mathbb{C}^8 . Then

$$X_{(4,8)} F_\bullet = \{H \in \text{Gr}(2, 8) \mid H \cap F_4 \neq \{0\}\},$$

as the remaining condition is that $H \subset F_8 = \mathbb{C}^8$, which always holds. The codimension of $X_{(4,8)} F_\bullet$ is $2(8-2) - (4-1) - (8-2) = 3$.

A *Schubert problem* is a list $\alpha = (\alpha^1, \dots, \alpha^s)$ with $\alpha^i \in \binom{[n]}{\ell}$ satisfying $|\alpha^1| + \dots + |\alpha^s| = \ell(n-\ell)$. By Kleiman's Transversality Theorem [16], if $F_\bullet^1, \dots, F_\bullet^s$ are general flags, then the intersection

$$(1.2) \quad X_{\alpha^1} F_\bullet^1 \cap X_{\alpha^2} F_\bullet^2 \cap \dots \cap X_{\alpha^s} F_\bullet^s,$$

is transverse and consists of finitely many points. Such an intersection (1.2) is an *instance* of the Schubert problem α .

Example 1.2. Suppose that $\alpha = ((4, 8), (4, 8), (4, 8), (4, 8))$, which is a Schubert problem on $\text{Gr}(2, 8)$. Let $F_\bullet^1, F_\bullet^2, F_\bullet^3, F_\bullet^4$ be general flags. The corresponding instance is

$$\begin{aligned} X_{(4,8)}F_\bullet^1 \cap X_{(4,8)}F_\bullet^3 \cap X_{(4,8)}F_\bullet^2 \cap X_{(4,8)}F_\bullet^4 \\ = \{H \in \text{Gr}(2, 8) \mid H \cap F_4^i \neq \{0\} \ i = 1, 2, 3, 4\}. \end{aligned}$$

That is, the 2-planes that have non-trivial intersection with four 4-planes in general position in \mathbb{C}^8 . This can be shown to have four solutions.

We may use the global Plücker coordinates to formulate a Schubert problem as the solutions to a system of equations. In the projective space $\mathbb{P}^{\binom{n}{\ell}-1}$, the Grassmannian is defined by quadratic Plücker equations and Schubert varieties are given by certain linear equations. This is not considered to be an efficient encoding of the Schubert problem.

For instance, in the Schubert problem of Example 1.2, Plücker space has dimension $\binom{8}{2} - 1 = 27$ and there are $\binom{8}{4} = 70$ linearly independent quadratic Plücker equations. These cut out the twelve-dimensional Grassmannian $\text{Gr}(2, 8)$ and each Schubert variety $X_{(4,8)}F_\bullet$ is cut out by six independent linear equations, for a total of 94 equations.

1.2. Local coordinates for Schubert varieties. Local coordinates for Schubert varieties in $\text{Gr}(\ell, n)$ are described in [8, Ch. 10]. Let M be a full rank $m \times n$ matrix with $m \leq n$. If $\ell \leq m$, let M_ℓ be the first ℓ rows of M and set $R_\ell(M)$ to be the row span of M_ℓ . For this, fix a basis $\mathbf{e}_1, \dots, \mathbf{e}_n$ of $V \simeq \mathbb{C}^n$ corresponding to the columns of M . If $a_\bullet = a_1 < \dots < a_t \subset [n]$ and $a_t \leq m$ then $R_{a_\bullet}(M)$ is the flag of linear subspaces

$$R_{a_1}(M) \subsetneq R_{a_2}(M) \subsetneq \dots \subsetneq R_{a_t}(M),$$

and we write $R_\bullet(M)$ for the complete flag obtained from an invertible $n \times n$ matrix M .

If N is an $n \times m$ matrix, we similarly have the column span $C_\ell(N)$ and the flag $C_{a_\bullet}(N)$. These are subspaces of the dual vector space V^* equipped with the basis $\mathbf{e}_1^*, \dots, \mathbf{e}_n^*$ of V^* dual to $\mathbf{e}_1, \dots, \mathbf{e}_n$, which corresponds to the rows of N .

The association $M \mapsto R_\ell(M)$ realizes the open subset of rank ℓ matrices in $\text{Mat}_{\ell \times n}(\mathbb{C})$ as a $GL(\ell, \mathbb{C})$ -principal bundle over the Grassmannian, called the Stiefel manifold. We identify subsets of the Stiefel manifold that parametrize Schubert varieties and their intersections, and call these subsets *Stiefel coordinates*.

Every ℓ -plane H is the row space of a unique echelon matrix, $M \in \text{Mat}_{\ell \times n}(\mathbb{C})$. That is, the last entry (the pivot) in row i of M is 1, and if α_i is the column of the pivot, then $\alpha_1 < \dots < \alpha_\ell$ and the 1 in position (i, α_i) is the only nonzero element of its column. Here are echelon matrices for $\alpha = (4, 8)$ and $\alpha = (2, 4, 5, 7)$ with $(\ell, n) = (2, 8)$ and $(4, 7)$, respectively, and where $m_{i,j}$ represents some complex number.

$$(1.3) \quad \begin{pmatrix} m_{1,1} & m_{1,2} & m_{1,3} & 1 & 0 & 0 & 0 & 0 \\ m_{2,1} & m_{2,2} & m_{2,3} & 0 & m_{2,5} & m_{2,6} & m_{2,7} & 1 \end{pmatrix} \quad \begin{pmatrix} m_{1,1} & 1 & 0 & 0 & 0 & 0 & 0 \\ m_{2,1} & 0 & m_{2,3} & 1 & 0 & 0 & 0 \\ m_{3,1} & 0 & m_{3,3} & 0 & 1 & 0 & 0 \\ m_{4,1} & 0 & m_{4,3} & 0 & 0 & m_{4,6} & 1 \end{pmatrix}$$

Let \mathcal{M}_α be the set of echelon matrices with pivots α . The number of unspecified entries in matrices in \mathcal{M}_α is $\dim(\alpha)$, which shows that $\mathcal{M}_\alpha \simeq \mathbb{C}^{\dim(\alpha)}$. Let $\mathbf{e}_1, \dots, \mathbf{e}_n$ be

an ordered basis of \mathbb{C}^n corresponding to the columns of our matrices. If we define the standard coordinate flag E_\bullet so that $E_j := \text{span}\{\mathbf{e}_1, \dots, \mathbf{e}_j\}$, then a ℓ -plane H has position α with respect to the flag E_\bullet , so that $H \in X_\alpha^\circ E_\bullet$, if and only if its echelon matrix has pivots α . As distinct echelon matrices give distinct ℓ -planes and vice-versa, this shows that $R_\ell: \mathcal{M}_\alpha \rightarrow X_\alpha^\circ E_\bullet$ is an isomorphism and thus $X_\alpha^\circ E_\bullet \simeq \mathbb{C}^{\dim(\alpha)}$. Since the Schubert cell is dense in the Schubert variety, \mathcal{M}_α gives local coordinates for the Schubert variety $X_\alpha E_\bullet$.

These echelon matrices \mathcal{M}_α also give local coordinates for arbitrary Schubert varieties. Let Φ be a full rank $n \times n$ matrix. If $M \in \mathcal{M}_\alpha$, then we invite the reader to check that $R_\ell(M\Phi) \in X_\alpha^\circ R_\bullet(\Phi)$.

Let E'_\bullet be the coordinate flag opposite to E_\bullet , where $E'_i := \text{span}\{\mathbf{e}_{n+1-i}, \dots, \mathbf{e}_n\}$ and $\dim E_i \cap E'_j = \max\{0, i+j-n\}$, so that E_\bullet and E'_\bullet are in linear general position. For $\beta \in \binom{[n]}{\ell}$, let \mathcal{M}^β be the set of echelon matrices in which row j has its first non-zero entry of 1 in column $n+1-\beta_{\ell+1-j}$, and this entry is the only nonzero entry in its column, but the remaining entries are unconstrained. Rotating a matrix 180° gives a bijection between \mathcal{M}^β and \mathcal{M}_β , and the map $M \mapsto R_\ell(M)$ is a bijection between \mathcal{M}^β and the Schubert cell $X_\beta^\circ E'_\bullet$. Here is a typical matrix in $\mathcal{M}^{(3,7,9)}$ for $\text{Gr}(3,9)$,

$$\begin{pmatrix} 1 & m_{1,2} & 0 & m_{1,4} & m_{1,5} & m_{1,6} & 0 & m_{1,8} & m_{1,9} \\ 0 & 0 & 1 & m_{2,4} & m_{2,5} & m_{2,6} & 0 & m_{2,8} & m_{2,9} \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & m_{3,8} & m_{3,9} \end{pmatrix}.$$

Suppose that we have indices $\alpha, \beta \in \binom{[n]}{\ell}$ with $\alpha_i + \beta_{\ell+1-i} \leq n$ for $i = 1, \dots, a$. Consider the set $\mathcal{M}_{\alpha, \beta}^\beta$ of matrices $(m_{i,j})$ where

$$\begin{aligned} m_{i,j} &= 0 & \text{if } i < \beta_{\ell+1-i} & \text{ or } \alpha_i < j \\ m_{i,j} &= 1 & \text{if } i = \beta_{\ell+1-i}, \end{aligned}$$

and $m_{i,j}$ is otherwise unconstrained. Here is a typical matrix in $M_{(6,8,10,11)}^{(5,7,8,11)}$,

$$\begin{pmatrix} 1 & m_{1,2} & m_{1,3} & m_{1,4} & m_{1,5} & m_{1,6} & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & m_{2,5} & m_{2,6} & m_{2,7} & m_{2,8} & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & m_{3,6} & m_{3,7} & m_{3,8} & m_{3,9} & m_{3,10} & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & m_{4,8} & m_{4,9} & m_{4,10} & m_{4,11} \end{pmatrix}.$$

If $M \in \mathcal{M}_{\alpha, \beta}^\beta$, then $R_\ell(M) \in X_\alpha E_\bullet \cap X_\beta E'_\bullet$. Indeed, as the last nonzero element m_{i, α_i} in row i is in column α_i , we conclude that $H \in X_\alpha E_\bullet$. To see that $H \in X_\beta E'_\bullet$, simply read the rows in reverse order.

A full rank $n \times n$ matrix Φ gives a complete flag $F_\bullet = R_\bullet(\Phi)$. If we let F'_i be the row span of the last i rows of Φ , we get a second flag $F'_\bullet = R_\bullet(w_0\Phi)$ which is opposite to F_\bullet in that $\dim F_i \cap F'_j = \max\{0, i+j-n\}$, and given any two such flags, there is a full rank matrix Φ which gives rise to them. We will refer to F_\bullet and F'_\bullet as the pair of opposite flags given by Φ . Note that E_\bullet and E'_\bullet are given by the identity matrix.

Proposition 1.3. *Let F_\bullet and F'_\bullet be opposite flags given by a full rank $n \times n$ matrix Φ , $F_\bullet = R_\bullet(\Phi)$ and $F'_\bullet = R_\bullet(w_0\Phi)$. The set \mathcal{M}_α of echelon matrices with pivots α is an affine space isomorphic to the Schubert cell $X_\alpha^\circ F_\bullet$ via the map $\mathcal{M}_\alpha \ni M \mapsto R_\ell(M\Phi)$.*

When $\alpha_i + \beta_{\ell+1-i} \leq n$ for $i = 1, \dots, \ell$, the set \mathcal{M}_α^β is an affine space which parametrizes a dense subset of the intersection $X_\alpha F_\bullet \cap X_\beta F'_\bullet$ via the map $\mathcal{M}_\alpha^\beta \ni M \mapsto R_\ell(MF)$. (If $\alpha_i + \beta_{\ell+1-i} > n$ for some i , then both \mathcal{M}_α^β and $X_\alpha F_\bullet \cap X_\beta F'_\bullet$ are empty.)

It only remains to show that \mathcal{M}_α^β parametrizes a dense subset of $X_\alpha F_\bullet \cap X_\beta F'_\bullet$. This follows as the map is injective on the locus where the entries m_{i,α_i} are simultaneously nonzero and the two varieties are irreducible of the same dimension $\ell(n-\ell) - |\alpha| - |\beta|$. For the intersection of Schubert varieties, this is because the flags are in general position, and for \mathcal{M}_α^β this is verified by counting parameters.

1.3. Determinantal equations for Schubert varieties. Let $\alpha \in \binom{[n]}{\ell}$ and let $F_\bullet = R_\bullet(\Phi)$ be a flag given by a full rank matrix Φ . The geometric conditions (1.1) defining a Schubert variety $X_\alpha F_\bullet$ give determinantal equations in local coordinates.

If $H = R_\ell(M)$, where $M \in \text{Mat}_{\ell \times n}(\mathbb{C})$, then $\dim(H \cap F_b) \geq c$ if and only if

$$\text{rank} \begin{pmatrix} M \\ \Phi_b \end{pmatrix} \leq \ell + b - c,$$

and so each of the square $\ell + b - c + 1$ submatrices (minors) of $\begin{pmatrix} M \\ \Phi_b \end{pmatrix}$ vanishes.

Proposition 1.4. *Let \mathcal{M} be a set of matrices parametrizing a subset Y of $\text{Gr}(\ell, \mathbb{C}^n)$, let Φ be a full-rank matrix, and let $\alpha \in \binom{[n]}{\ell}$. Then, in the coordinates $M \in \mathcal{M}$ for Y , the intersection $Y \cap X_\alpha R_\bullet(\Phi)$ of Y with the Schubert variety is defined by the vanishing of the square $\ell + \alpha_i - i + 1$ minors of $\begin{pmatrix} M \\ \Phi_{\alpha_i} \end{pmatrix}$, for $i = 1, \dots, \ell$.*

This set $I(\alpha, \Phi)$ of minors defining $X_\alpha F_\bullet$ is typically more than is necessary. For $\alpha \in \binom{[n]}{\ell}$, let $\|\alpha\|$ be the number of sequences $\beta \in \binom{[n]}{\ell}$ such that $\beta \not\leq \alpha$.

Lemma 1.5. *The ideal generated by $I(\alpha, \Phi)$ that defines $Y \cap X_\alpha F_\bullet$ is generated by at most $\|\alpha\|$ polynomials that are linear combinations of full rank minors of M .*

Proof. In the Plücker embedding of $\text{Gr}(\ell, n)$, the Plücker coordinates $\{p_\beta \mid \beta \in \binom{[n]}{\ell}\}$ form a basis of the linear forms. If M is a $\ell \times n$ matrix with $H = R_\ell(M)$, then $p_\beta(H)$ is the $\ell \times \ell$ minor of the matrix M formed by its columns indexed by β .

The Schubert variety $X_\alpha E_\bullet$ given by the coordinate flag E_\bullet is defined in $\text{Gr}(\ell, n)$ by the vanishing of the Plücker coordinates p_β for $\beta \not\leq \alpha$. If g is a matrix sending E_\bullet to F_\bullet , then it acts on Plücker space sending $X_\alpha E_\bullet$ to $X_\alpha F_\bullet$ and thus $X_\alpha F_\bullet$ is defined by the linear forms $g(p_\beta)$ with $\beta \not\leq \alpha$. Thus $Y \cap X_\alpha F_\bullet$ is defined by the linear combinations of minors $\{g(p_\beta)(H) \mid \beta \not\leq \alpha\}$, and this set has cardinality $\|\alpha\|$. \square

Corollary 1.6. *The Schubert variety $X_\alpha F_\bullet$ is defined by $|\alpha|$ equations if and only if either $\ell = 1$, $\ell = n - 1$, or $\alpha = (n - \ell, n - \ell + 2, \dots, n)$, so that $1 = |\alpha| = \|\alpha\|$.*

Proof. By Lemma 1.5, we must describe when $|\alpha| = \|\alpha\|$. If $\ell = 1$ or $\ell = n - 1$, then $\text{Gr}(\ell, n) \simeq \mathbb{P}^{n-1}$, Schubert varieties are linear subspaces, and so $|\alpha| = \|\alpha\|$. If $\alpha =$

$(n-\ell, n-\ell+2, \dots, n)$ then $X_\alpha F_\bullet$ is a hypersurface Schubert variety, defined by the single equation $\det \begin{pmatrix} M \\ F_{n-\ell} \end{pmatrix} = 0$. In all other cases, $|\alpha| < \|\alpha\|$. \square

Write \square for the Schubert condition $(n-\ell, n-\ell+2, \dots, n)$.

Example 1.7. The set $\mathcal{M}_{(7,8)}$ of matrices of the form $M = [X : I_2]$ for $X \in \text{Mat}_{2 \times 6}(\mathbb{C})$ parametrizes $X_{(7,8)}^\circ E_\bullet$, which is a dense subset of the Grassmannian $\text{Gr}(2, 8)$. In this cell, the Schubert variety $X_{(4,8)} F_\bullet$ is defined by the vanishing of the $\binom{8}{6} = 28$ minors of size 6×6 of $\begin{pmatrix} M \\ \Phi_4 \end{pmatrix}$ (when $i = 1$), and there are no equations for $i = 2$. These minors are dependent. For Φ general, these 28 minors span a six-dimensional linear space of polynomials, by Lemma 1.5. Indeed, the set of sequences $\{\beta \in \binom{[8]}{2} \mid (4, 8) \not\preceq \beta\}$ is

$$\{(7, 8), (6, 8), (5, 8), (6, 7), (5, 7), (5, 6)\}.$$

Let $(\alpha^1, \dots, \alpha^s)$ be a Schubert problem on $\text{Gr}(\ell, n)$ and suppose that $F_\bullet^1, \dots, F_\bullet^s$ are flags in general position, giving an instance (1.2) of this Schubert problem. We use Propositions 1.4 and 1.3 to give three formulations of this instance in terms of equations and local coordinates. Let Φ^1, \dots, Φ^s be full-rank matrices giving the flags, $F_\bullet^i := R_\bullet(\Phi^i)$.

1.3.1. *Local coordinates for $\text{Gr}(\ell, n)$.* Let $\mathcal{M} := \mathcal{M}_{(n-\ell+1, \dots, n)}$ be the set of matrices $(X : I_\ell)$ where $X \in \text{Mat}_{\ell \times (n-\ell)}(\mathbb{C})$, which parametrizes the dense Schubert cell $X_{(n-\ell+1, \dots, n)}^\circ E_\bullet$. Then the instance (1.2) is defined in \mathcal{M} by the collection of ideals $I(\alpha^i, \Phi^i)$, for $i = 1, \dots, s$. This uses $\ell(n-\ell)$ variables.

1.3.2. *Local coordinates for $X_{\alpha^s} F_\bullet^s$.* Assume that the coordinates for \mathbb{C}^n have been chosen so that $F_\bullet^s = E_\bullet$, the standard coordinate flag—this may be achieved by replacing each matrix Φ^i by $\Phi^i(\Phi^s)^{-1}$. Let $\mathcal{M} := \mathcal{M}_{\alpha^s}$ be the set of matrices parametrizing the Schubert cell $X_{\alpha^s}^\circ E_\bullet$. Then the instance (1.2) is defined in \mathcal{M} by the collection of ideals $I(\alpha^i, \Phi^i)$, for $i = 1, \dots, s-1$. This uses $\ell(n-\ell) - |\alpha^s|$ variables.

1.3.3. *Local coordinates for $X_{\alpha^s} F_\bullet^s \cap X_{\alpha^{s-1}} F_\bullet^{s-1}$.* Assume that the coordinates for \mathbb{C}^n have been chosen so that $F_\bullet^s = E_\bullet$ and $F_\bullet^{s-1} = E'_\bullet$ are the standard coordinate flags. This may be achieved by choosing \mathbf{e}_i to be any nonzero vector in the one-dimensional linear subspace $F_i^s \cap F_{n+1-i}^{s-1}$ for each $i = 1, \dots, n$. Then the instance (1.2) is defined in $\mathcal{M} = \mathcal{M}_{\alpha^s}^{\alpha^{s-1}}$ by the collection of ideals $I(\alpha^i, \Phi^i)$, for $i = 1, \dots, s-2$. This uses $\ell(n-\ell) - |\alpha^s| - |\alpha^{s-1}|$ variables.

Remark 1.8. In each of these formulations, if the flags F_\bullet^i are general then the set of solutions to the system of equations in \mathcal{M} will give the set of solutions to the Schubert problem, under the map $\mathcal{M} \ni M \mapsto R_\ell(M)$. Indeed, the parametrization map $\mathcal{M} \rightarrow \text{Gr}(\ell, n)$ is one-to-one on an open subset of \mathcal{M} and the image is dense in the corresponding Schubert variety or intersection of Schubert varieties. By Kleiman's Transversality Theorem, there will be no solutions in the lower dimensional subvariety where the map is not injective.

2. PRIMAL-DUAL FORMULATION OF SCHUBERT PROBLEMS

Associating a ℓ -plane H in $V = \mathbb{C}^n$ to its annihilator H^\perp in the linear dual V^* of V identifies the Grassmannian $\text{Gr}(\ell, V)$ with $\text{Gr}(n-\ell, V^*)$. This identification sends Schubert varieties to Schubert varieties. The graph of this map is defined by bilinear equations in

Stiefel coordinates. We explain how the parametrization of Schubert varieties and a twist on the classical reduction to the diagonal leads to a formulation of any Schubert problem on $\text{Gr}(\ell, V)$ as a square system of bilinear equations.

2.1. Parametrizing dual Schubert varieties. Let $\mathbf{e}_1, \dots, \mathbf{e}_n$ be a basis for the vector space V and $\mathbf{e}_1^*, \dots, \mathbf{e}_n^*$ be the dual basis for V^* . Write \perp for the canonical map $\perp: H \rightarrow H^\perp$ between $\text{Gr}(\ell, V)$ and $\text{Gr}(n-\ell, V^*)$. We first identify the image $\perp(X_\alpha F_\bullet)$ of a Schubert variety under the map \perp .

Each flag F_\bullet on V has a corresponding *dual flag* F_\bullet^\perp on V^* , where $(F_\bullet^\perp)_i = (F_{n-i})^\perp$,

$$F_\bullet^\perp : (F_{n-1})^\perp \subsetneq (F_{n-2})^\perp \subsetneq \cdots \subsetneq (F_1)^\perp \subsetneq V^*,$$

For $\alpha \in \binom{[n]}{\ell}$, define the sequence $\alpha^\perp \in \binom{[n]}{n-\ell}$ by $j \in \alpha^\perp$ if and only if $n+1-j \notin \alpha$. For example, if $\ell = 4$ and $n = 7$, then $(2, 4, 5, 7)^\perp = (2, 5, 7)$. Note that

$$[n] = \alpha^\perp \cup \{n+1-\alpha_i \mid i = 1, \dots, \ell\},$$

and thus $n-i = \#(\alpha \cap \{i+1, \dots, n\}) + \#(\alpha^\perp \cap [n-i])$. (Here, $\#S$ is the cardinality of the set S , and $(\alpha \cap S) := (\alpha_i \mid \alpha_i \in S)$.)

Lemma 2.1. *For a Schubert variety $X_\alpha F_\bullet \subset \text{Gr}(\ell, V)$, we have $\perp(X_\alpha F_\bullet) = X_{\alpha^\perp} F_\bullet^\perp$.*

Note that $X_\alpha F_\bullet = \perp(X_{\alpha^\perp} F_\bullet^\perp)$. We call $X_\alpha F_\bullet$ and $X_{\alpha^\perp} F_\bullet^\perp$ *dual Schubert varieties*.

Proof. Observe that if F_\bullet is a flag and H a linear subspace, then $\dim H \cap F_b \geq a$ implies that $\dim H \cap F_{b+1} \geq a$. Thus the definition (1.1) of Schubert variety is equivalent to

$$(2.1) \quad X_\alpha F_\bullet := \{H \in \text{Gr}(\ell, V) \mid \dim(H \cap F_i) \geq \#(\alpha \cap [i]), \text{ for } i = 1, \dots, n\}.$$

For every $H \in \text{Gr}(\ell, V)$ and all $i = 1, \dots, n$, the following are equivalent.

$$\begin{aligned} \dim H \cap F_i &\geq \#(\alpha \cap [i]) \\ &\Leftrightarrow \dim(\text{span}\{H, F_i\}) \leq \ell+i - \#(\alpha \cap [i]) = i + \#(\alpha \cap \{i+1, \dots, n\}) \\ &\Leftrightarrow \dim(\text{span}\{H, F_i\}^\perp) \geq n-i - \#(\alpha \cap \{i+1, \dots, n\}) \\ &\Leftrightarrow \dim(H^\perp \cap F_{n-i}^\perp) \geq n-i - \#(\alpha \cap \{i+1, \dots, n\}). \end{aligned}$$

Since $n-i - \#(\alpha \cap \{i+1, \dots, n\}) = \#(\alpha^\perp \cap [n-i])$, the lemma follows from (1.1). \square

Observe that we have

$$(2.2) \quad \begin{pmatrix} m_{1,1} & 1 & 0 & 0 & 0 & 0 & 0 \\ m_{2,1} & 0 & m_{2,3} & 1 & 0 & 0 & 0 \\ m_{3,1} & 0 & m_{3,3} & 0 & 1 & 0 & 0 \\ m_{4,1} & 0 & m_{4,3} & 0 & 0 & m_{4,6} & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 \\ -m_{1,1} & 0 & 0 \\ 0 & 1 & 0 \\ -m_{2,1} & -m_{2,3} & 0 \\ -m_{3,1} & -m_{3,3} & 0 \\ 0 & 0 & 1 \\ -m_{4,1} & -m_{4,3} & -m_{4,6} \end{pmatrix} = \begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}.$$

Thus if $H \in X_\alpha E_\bullet$ is the row space of the 4×7 matrix $M \in \mathcal{M}_{(2,4,5,7)}$ in (2.2), then $H^\perp \in X_{(2,4,5,7)^\perp} E_\bullet^\perp$ is the column space of the 7×3 matrix in (2.2) and vice-versa. Here, the dual basis $\mathbf{e}_1^*, \dots, \mathbf{e}_n^*$ corresponds to the rows of the 7×3 matrix. In particular,

$\mathcal{M}^{(2,5,7)}$ parametrizes $X_{(2,4,5,7)\perp}E_{\bullet}^{\perp}$ via the map $M \mapsto C_3(M^T)$. Here, we write M^T for the transpose of a matrix M .

Lemma 2.2. *The Schubert cell $X_{\alpha\perp}E_{\bullet}^{\perp}$ is parametrized by matrices in $\mathcal{M}^{\alpha\perp}$ via $M \mapsto C_{n-\ell}(M^T)$. Similarly, the intersection $X_{\alpha\perp}E_{\bullet}^{\perp} \cap X_{\beta\perp}E_{\bullet}^{\perp}$ is parametrized by matrices in $\mathcal{M}_{\beta\perp}^{\alpha\perp}$ via $M \mapsto C_{n-\ell}(M^T)$.*

We leave the proof of this lemma to the reader.

Note that if $H \in \mathcal{M}_{\alpha}$ and $K \in \mathcal{M}^{\alpha\perp}$, then each of the equations $HK^T = 0_{\ell \times (n-\ell)}$ is either trivial or of the form $m_{i,j} + k_{t,j} = 0$ as in (2.2). Define \mathcal{N}_{α} to be the set of $n \times (n-\ell)$ matrices $\{M^T \mid M \in \mathcal{M}^{\alpha\perp}\}$. Then \mathcal{N}_{α} parametrizes $X_{\alpha\perp}E_{\bullet}^{\perp}$ via $N \mapsto C_{n-\ell}(N)$. As $\Phi\Phi^{-1} = I$, if $F_{\bullet} = R_{\bullet}(\Phi)$ then $(F_{\bullet})^{\perp}$ is the flag whose i plane is the span of the last i columns of Φ^{-1} . Thus \mathcal{N}_{α} parametrizes $X_{\alpha\perp}E_{\bullet}^{\perp}$ via $N \mapsto C_{n-\ell}(\Phi^{-1}N)$.

2.2. The primal-dual formulation of a Schubert problem. Let $\Delta: \text{Gr}(\ell, V) \rightarrow \text{Gr}(\ell, V) \times \text{Gr}(n-\ell, V^*)$ be the graph of the isomorphism $\perp: \text{Gr}(\ell, V) \rightarrow \text{Gr}(n-\ell, V^*)$. We call Δ the *dual diagonal* map and will identify $\text{Gr}(\ell, V)$ with its image under Δ . In this context, the classical reduction to the diagonal becomes the following.

Lemma 2.3. *Let $A, B \subset \text{Gr}(\ell, V)$. Then $\Delta(A \cap B) = (A \times \perp(B)) \cap \Delta(\text{Gr}(\ell, V))$.*

Proof. Let $H \in A \cap B$. Then $\Delta(H) = (H, H^{\perp}) \in A \times \perp(B)$, which establishes the containment \subset . For the other containment, let $(H, K) \in (A \times \perp(B)) \cap \Delta(\text{Gr}(\ell, V))$. Then $H \in A$ and $H^{\perp} = K \in \perp(B)$, so that $H \in B$, which completes the proof. \square

We use Lemma 2.3 to express a Schubert problem as a complete intersection given by bilinear equations. Suppose that M is a $\ell \times n$ matrix whose row space H is a ℓ -plane in V and N is a $n \times (n-\ell)$ matrix whose column space K is a $(n-\ell)$ -plane in V^* . (The coordinates of the matrices—columns for M and rows for N —are with respect to the bases \mathbf{e}_i and \mathbf{e}_j^* , respectively.) Then $H^{\perp} = K$ if and only if $MN = 0_{\ell \times (n-\ell)}$, giving $\ell(n-\ell)$ bilinear equations in the entries of M and N for $\Delta(\text{Gr}(\ell, V))$. We record this fact.

Lemma 2.4. *Let A, B be two subsets of $\text{Gr}(\ell, V)$ and suppose that \mathcal{M} is a set of $\ell \times n$ matrices parametrizing A (via row span) and that \mathcal{N} is a set of $n \times (n-\ell)$ matrices parametrizing $\perp(B)$ (via column span). Then the subset of $\mathcal{M} \times \mathcal{N}$ defined by the equations $MN = 0_{\ell \times (n-\ell)}$ parametrizes $\Delta(A \cap B)$ as a subset of $A \times \perp(B)$.*

We deduce the primal-dual formulation of the Schubert variety $X_{\alpha}F_{\bullet}$.

Corollary 2.5. *Suppose that $F_{\bullet} = R_{\bullet}(\Phi)$ is a general flag. Then the bilinear equations $M\Phi^{-1}N = 0_{k \times (n-\ell)}$ on pairs $(M, N) \in \mathcal{M}_{(n-\ell+1, \dots, n)} \times \mathcal{N}_{\alpha}$ define $\Delta(X_{(n-\ell+1, \dots, n)}^{\circ}E_{\bullet} \cap X_{\alpha}^{\circ}F_{\bullet})$ as a system of $\ell(n-\ell)$ equations in $2\ell(n-\ell) - |\alpha|$ variables.*

We require F_{\bullet} to be sufficiently general so that $X_{(n-\ell+1, \dots, n)}^{\circ}E_{\bullet} \cap X_{\alpha}^{\circ}F_{\bullet}$ is nonempty, for then it is a dense subset of $X_{\alpha}F_{\bullet}$ (as $X_{(n-\ell+1, \dots, n)}^{\circ}E_{\bullet}$ is a dense subset of $\text{Gr}(\ell, V)$). Since $\dim \Delta(X_{(n-\ell+1, \dots, n)}^{\circ}E_{\bullet} \cap X_{\alpha}^{\circ}F_{\bullet}) = \dim(X_{\alpha}F_{\bullet}) = \ell(n-\ell) - |\alpha|$, this system of equations and variables exhibits $X_{\alpha}F_{\bullet}$ as a complete intersection of bilinear equations in local coordinates which we call the *primal-dual formulation* of the Schubert variety $X_{\alpha}F_{\bullet}$.

2.3. Primal-dual formulation of a Schubert problem. Extending this primal-dual formulation to Schubert problems uses a dual diagonal map to the small diagonal in a larger product of Grassmannians. Define

$$\Delta^s : \text{Gr}(\ell, V) \rightarrow \text{Gr}(\ell, V) \times (\text{Gr}(n - \ell, V^*))^{s-1},$$

by sending $H \mapsto (H, H^\perp, \dots, H^\perp)$.

Lemma 2.6. *Let $A_1, \dots, A_s \subset \text{Gr}(\ell, V)$. Then*

$$\Delta^s(A_1 \cap \dots \cap A_s) = (A_1 \times \perp(A_2) \times \dots \times \perp(A_s)) \cap \Delta^s(\text{Gr}(\ell, V)).$$

Lemma 2.4 extends to the dual diagonal of many factors.

Lemma 2.7. *Let $A_1, \dots, A_s \subset \text{Gr}(\ell, V)$, suppose that \mathcal{M} is a set of $\ell \times n$ matrices parametrizing A_1 and \mathcal{N}_i is set of $n \times (n - \ell)$ matrices parametrizing $\perp(A_i)$, for each $i = 2, \dots, s$. Then $\Delta^s(A_1 \cap \dots \cap A_s)$ is the subset of $A_1 \times \perp(A_2) \times \dots \times \perp(A_s)$ defined in the parametrization $\mathcal{M} \times \mathcal{N}_2 \times \dots \times \mathcal{N}_s$ by the equations $MN_i = 0_{\ell \times (n - \ell)}$ for $i = 2, \dots, s$, where $(M, N_2, \dots, N_s) \in \mathcal{M} \times \mathcal{N}_2 \times \dots \times \mathcal{N}_s$.*

We deduce our main theorem of this section.

Theorem 2.8. *Let $(\alpha^1, \dots, \alpha^s)$ be a Schubert problem for $\text{Gr}(\ell, V)$ and suppose that $F_\bullet^1, \dots, F_\bullet^s$ are general flags given by matrices Φ^1, \dots, Φ^s . Then the system of equations*

$$(2.3) \quad (M\Phi^s)(\Phi^i)^{-1}N_i = 0_{\ell \times (n - \ell)} \quad \text{for } i = 1, \dots, s-1,$$

where $(M, N_1, \dots, N_{s-1}) \in M_{\alpha^s} \times N_{\alpha^1} \times \dots \times N_{\alpha^{s-1}}$ defines the instance of the Schubert problem (1.2) as the common zeroes of $(s-1) \cdot \ell(n - \ell)$ bilinear equations in

$$\sum_{i=1}^s (\ell(n - \ell) - |\alpha^i|) = (s-1) \cdot \ell(n - \ell)$$

variables, exhibiting the Schubert problem as a square system of equations.

Proof. A sufficiently general instance of the Schubert problem is zero-dimensional. The result follows from Lemma 2.7 in which A_1 is the open subset of $X_{\alpha^s} F_\bullet^s$ parametrized by \mathcal{M}_{α^s} and for $i < s$, A_i is the open subset of $X_{\alpha^i} F_\bullet^i$ with $\perp(A_i)$ parametrized by $\mathcal{M}_{\alpha^i \perp}$. This gives $(s-1) \cdot \ell(n - \ell)$ bilinear equations in $(s-1) \cdot \ell(n - \ell)$ variables, and the generality of the flags ensures that all points in the intersection of Schubert varieties (1.2) lie in the intersection of Schubert cells. \square

Remark 2.9. Theorem 2.8 provides a formulation of an instance of a Schubert problem as a square system, to which the certification afforded by Smale's α -theory may be applied. This rectifies the fundamental obstruction to using numerical methods in place of certified symbolic methods for solving Schubert problems.

Remark 2.10. The bilinear equations (2.3) are partitioned into $s-1$ blocks of $\ell(n - \ell)$ equations. Each block of equations is linear in the $\ell(n - \ell) - |\alpha^s|$ variables in M_{α^s} , with the i th block also linear in the variables in \mathcal{N}_{α^i} , which occur in no other block. This bilinear structure can be exploited in the certification via α -theory. Since all equations

are quadratic, the supremum defining the γ term for α -theory is taken over only one term. One can either explicitly compute this term or could utilize the bilinear structure to produce an upper bound that is sharper than, say, Proposition 3 from § I-3 of [27]. Similar bilinear structure was exploited for eigenvalue and generalized eigenvalue problems in [1, 6].

3. IMPROVEMENTS TO THE PRIMAL-DUAL FORMULATION

The formulation of a Schubert problem in Theorem 2.8 solves the problem of certifiability, but is not particularly efficient. For example, the Schubert problem of Example 1.2 involves $3 \cdot 12 = 36$ equations and variables. We present two improvements to Theorem 2.8. Let $\alpha = (\alpha^1, \dots, \alpha^s)$ be a Schubert problem for $\text{Gr}(\ell, V)$ and $F_\bullet^1, \dots, F_\bullet^s$ be general flags in V .

3.1. Efficient local coordinates. The formulation of Theorem 2.8 may be improved using the local coordinates \mathcal{M}_α^β for $X_\alpha F_\bullet \cap X_\beta F'_\bullet$. This reduces the number of equations/variables by $\lfloor \frac{s}{2} \rfloor \ell(n - \ell)$.

Adding a trivial condition $(n - \ell + 1, \dots, n)$ to α if necessary, we will assume that $s = 2k$. For each $i = 1, \dots, k$, let Φ^i be a matrix such that $F_\bullet^{2i-1} = R_\bullet(\Phi^i)$ and $F_\bullet^{2i} = R_\bullet(w_0 \Phi^i)$. This requires that F_\bullet^{2i-1} and F_\bullet^{2i} are in linear general position. For $\alpha, \beta \in \binom{[n]}{\ell}$ with $\alpha_i + \beta_{\ell+1-i} \leq n$ for all i , let \mathcal{N}_α^β be the set of $n \times (n - \ell)$ matrices $\{M^T \mid M \in \mathcal{M}_{\beta^\perp}^{\alpha^\perp}\}$.

Theorem 3.1. *The system of equations*

$$(3.1) \quad (M\Phi^k)(\Phi^i)^{-1}N_i = 0_{k \times (n-\ell)} \quad \text{for } i = 1, \dots, k-1,$$

where $(M, N_1, \dots, N_{k-1}) \in \mathcal{M}_{\alpha_{2k-1}^{\alpha^{2k}}} \times \mathcal{N}_{\alpha^1}^{\alpha^2} \times \dots \times \mathcal{N}_{\alpha_{2k-3}^{\alpha^{2k-2}}}$ defines the instance of the Schubert problem (1.2) as the common zeroes of $(k-1) \cdot \ell(n-\ell)$ bilinear equations in

$$\sum_{i=1}^k (\ell(n-\ell) - |\alpha^{2i-1}| - |\alpha^{2i}|) = (k-1) \cdot \ell(n-\ell)$$

variables, exhibiting the Schubert problem as a square system of equations.

We omit the proof, as it is similar to that of Theorem 2.8.

Remark 3.2. The equations (3.1) exhibit a block structure similar to the equations (2.3), and the same comments made in Remark 2.10 about their structure also apply here.

We determine the reduction in the numbers of equations/variables obtained by using the formulation of Theorem 3.1 in place of the formulation of Theorem 2.8. Set $\nu := \ell(n - \ell)$. If $s = 2k$ is even, then we do not add a trivial Schubert condition to α . In this case the formulation of Theorem 2.8 involves $(2k-1) \cdot \nu$ equations while that of Theorem 3.1 involves $(k-1) \cdot \nu$, a reduction of $k\nu$.

If $s = 2k-1$ is odd, then we add a trivial Schubert condition and the number of equations from Theorem 2.8 is $(2k-2) \cdot \nu$ and $(k-1) \cdot \nu$ in Theorem 3.1, for a reduction of $(k-1) \cdot \nu$. Thus the formulation of Theorem 3.1 reduces the numbers of equations/variables by $\lfloor \frac{s}{2} \rfloor \cdot \ell(n - \ell)$ from that of Theorem 2.8.

3.2. Exploiting hypersurface Schubert varieties. A further improvement is possible when the Schubert problem α includes some codimension 1 conditions, \square . Recall that the Schubert variety $X_{\square}F_{\bullet}$ is defined by a single determinant. Suppose that α has the form $(\alpha^1, \dots, \alpha^{2k}, \square, \dots, \square)$, where there are at least t occurrences of \square and $s = 2k + t$.

For each $i = 1, \dots, k$, let Φ^i be a matrix giving the pair of opposite flags $F_{\bullet}^{2i-1} = R_{\bullet}(\Phi)$ and $F_{\bullet}^{2i} = R_{\bullet}(w_0\Phi)$ as in Subsection 3.1. For each $j = 1, \dots, t$, let Ψ^j be a $(n-\ell) \times n$ matrix with $R_{n-\ell}(\Psi^j)$ the linear subspace $F_{n-\ell}^{2k+j}$ in the flag F_{\bullet}^{2k+j} .

Theorem 3.3. *With these definitions, the system of equations*

$$(3.2) \quad \begin{aligned} (M\Phi^k)(\Phi^i)^{-1}N_i &= 0_{\ell \times (n-\ell)} \quad \text{for } i = 1, \dots, k-1, \\ \det \begin{pmatrix} M \\ \Psi^j \end{pmatrix} &= 0 \quad \text{for } j = 1, \dots, t, \end{aligned}$$

where $(M, N_1, \dots, N_{k-1}) \in \mathcal{M}_{\alpha_{2k-1}^{2k}} \times \mathcal{N}_{\alpha_1^2} \times \dots \times \mathcal{N}_{\alpha_{2k-3}^{2k-2}}$ defines the instance of the Schubert problem (1.2) as the common zeroes of $(k-1) \cdot \ell(n-\ell)$ bilinear equations and t determinantal equations in

$$\sum_{i=1}^k (\ell(n-\ell) - |\alpha^{2i-1}| - |\alpha^{2i}|) = k \cdot \ell(n-\ell) - \sum_{i=1}^{2k} |\alpha^i| = (k-1) \cdot \ell(n-\ell) + t$$

variables, exhibiting the Schubert problem as a square system of equations.

Remark 3.4. The previous comments based on the block structure of equations remain valid for the formulation (3.2).

Suppose now that α has a conditions $\alpha^1, \dots, \alpha^a$ with $|\alpha^i| > 1$ for $i = 1, \dots, a$ and $b \geq 1$ occurrences of the codimension 1 condition \square . If $a = 2k$ is even, then we apply Theorem 3.3 with $t = b$. If $a = 2k-1$ is odd, then we set $\alpha^{2k} = \square$ and apply Theorem 3.3 with $t = b-1$ (necessarily, $b \geq 2$ for this). Table 1 shows the gain in efficiency from using Theorem 3.3 in place of Theorem 3.1 when we have hypersurface Schubert conditions, recording the reduction in the number of equations/variables. Set $\nu := \ell(n-\ell)$. In all cases we have a net reduction.

TABLE 1. Efficiency gain with hypersurface Schubert conditions

$(a, b) \pmod 2$	$(0, 0)$	$(0, 1)$	$(1, 0)$	$(1, 1)$
reduction	$\lfloor \frac{b}{2} \rfloor \nu - b$	$\lfloor \frac{b+1}{2} \rfloor \nu - b$	$\lfloor \frac{b}{2} \rfloor \nu - b + 1$	$\lfloor \frac{b}{2} \rfloor \nu - b + 1$

Example 3.5. We compare all these formulations for a Schubert problem in $\text{Gr}(3, 9)$. Let $\beta := (4, 8, 9)$, which has $|\beta| = 3$ and corresponds to the condition that a 3-plane H meets a fixed 4-plane nontrivially. Set $\alpha := (\beta, \beta, \beta, \beta, \square, \square, \square, \square, \square)$, which has 437 solutions.

The Grassmannian $\text{Gr}(3, 9)$ is defined in Plücker space \mathbb{P}^{83} by 1050 independent quadratic Plücker equations. Since $\|\beta\| = 10$ and $\|\square\| = 1$, any instance of the Schubert problem α is defined in Plücker space by 1050 quadratic and 46 linear equations.

The formulation of § 1.3.1 in local coordinates $(X: I_a)$ for $\text{Gr}(3, 9)$ with $X \in \text{Mat}_{3 \times 6}$ uses 18 variables and 46 independent cubic minors/determinants to express the Schubert problem α . Using the local coordinates \mathcal{M}_β as in § 1.3.2, reduces this to 15 variables and 36 cubic equations, while the formulation of § 1.3.3 using local coordinates \mathcal{M}_β^β involves 12 variables and 26 cubic equations. These formulations are all overdetermined.

In contrast, the formulation of Theorem 2.8 for α uses $9 \cdot 18 = 162$ variables and 162 bilinear equations, that of Theorem 3.1 uses $4 \cdot 18 = 72$ variables and bilinear equations, while that of Theorem 3.3 uses 24 variables with 18 bilinear equations and six cubic determinants, coming from the six codimension one conditions \square .

To test the square formulation of 24 variables, we used Bertini [2] to solve an instance of this Schubert problem given by random real flags. The computation consumed approximately 20.37 gigaHertz-hours of processing power to calculate 437 approximate solutions, and the output suggests that 85 solutions are real 3-planes while the other 352 are nonreal. For this, we used regeneration which is not specific to Schubert calculus and which may be done efficiently in parallel.

We used rational arithmetic in alphaCertified [13] to prove that the 437 solutions given by Bertini are approximate solutions to the computed instance, and that they correspond to distinct solutions. Since they are distinct, theorems of Schubert calculus guarantee that we have approximations for every solution to the instance. We also verified that 85 of the approximate solutions correspond to real solutions. This calculation took approximately 2.00 gigaHertz-hours of processing power. Certification may be done in parallel, efficiently using up to 437 processors.

4. PRIMAL-DUAL FORMULATION FOR FLAG MANIFOLDS

A flag manifold is the set of all flags of linear subspaces of specified dimensions, and is therefore a generalization of the Grassmannian. Flag manifolds have Schubert varieties and Schubert problems, and these may be formulated as systems of determinantal equations in local coordinates. Like the Grassmannian, there is a primal-dual formulation of Schubert problems on flag manifolds.

Let V be an n -dimensional complex vector space and let $a_\bullet: 0 < a_1 < \dots < a_t < n$ be a sequence of integers. A *flag of type a_\bullet* is an increasing sequence

$$V_{a_\bullet} : V_{a_1} \subsetneq V_{a_2} \subsetneq \dots \subsetneq V_{a_t} \subsetneq V,$$

of linear subspaces of V with $\dim V_{a_i} = a_i$ for $i = 1, \dots, t$. A flag is *complete* if $a_\bullet = [n-1]$. The set of all flags of type a_\bullet forms the flag manifold $\mathbb{F}\ell(a_\bullet, V)$, which is a homogeneous space for the general linear group $\text{GL}(V)$. If we set $a_0 := 0$, then it has dimension

$$\dim a_\bullet := \sum_{i=1}^t (a_i - a_{i-1})(n - a_i).$$

If P_{a_\bullet} is the subgroup stabilizing a flag of type a_\bullet , then $\mathbb{F}\ell(a_\bullet, V) \simeq \text{GL}(n, \mathbb{C})/P_{a_\bullet}$. The Weyl group W_{a_\bullet} of P_{a_\bullet} is the Young subgroup $S_{a_1} \times S_{a_2 - a_1} \times \dots \times S_{n - a_t}$ of the symmetric group S_n , which is the Weyl group of $\text{GL}(n, \mathbb{C})$.

Given a complete flag F_\bullet , the flag manifold has a decomposition into Schubert cells whose closures are Schubert varieties. A Schubert cell consists of all flags V_{a_\bullet} having the same specified position with respect to the reference flag F_\bullet . These positions are indexed by permutations $w \in S_n$ on n letters with descents in the set a_\bullet ,

$$W^{a_\bullet} := \{w \in S_n \mid w(i) < w(i+1) \text{ if } i \notin a_\bullet\},$$

which are minimal length representatives of cosets of W_{a_\bullet} in S_n . The Schubert variety associated to a (complete) flag F_\bullet and permutation $w \in W^{a_\bullet}$ is

$$(4.1) \quad X_w F_\bullet := \{V_{a_\bullet} \in \mathbb{F}\ell(a_\bullet, V) \mid \dim V_{a_i} \cap F_j \geq \#\{k \leq a_i \mid w(k) \leq j\}\}.$$

For $a \in a_\bullet$ we have the map $w \mapsto w|_a := \{w(1), \dots, w(a)\}$ which sends W^{a_\bullet} to $\binom{[n]}{a}$. An equivalent form of the definition (4.1) is that

$$(4.2) \quad V_{a_\bullet} \in X_w F_\bullet \text{ if and only if } V_a \in X_{w|_a} F_\bullet \subset \text{Gr}(a, V) \text{ for all } a \in a_\bullet.$$

When $t = 1$ so that $a_\bullet = \{a\}$, then $\mathbb{F}\ell(a_\bullet, V) = \text{Gr}(a, V)$, permutations in $W^{\{a\}}$ are in bijection with $\binom{[n]}{a}$ via $w \mapsto w|_a$, and we obtain the same Schubert varieties as before.

As explained in [8, Ch. 10], matrices in partial echelon form parametrize Schubert cells. For $w \in W^{a_\bullet}$ let \mathcal{M}_w be the set of those $a_t \times n$ matrices $(m_{i,j})$ where for all i, j ,

$$\begin{aligned} m_{i,j} &= 1 && \text{if } j = w(i) \\ m_{i,j} &= 0 && \text{if } j > w(i) \text{ or } j = w(k) \text{ for some } k < i. \end{aligned}$$

Note then that $m_{i,j}$ is 1 or 0 unless $j = w(k) < w(i)$ for some $k > i$. Thus $\mathcal{M}_w \simeq \mathbb{C}^{\ell(w)}$, where $\ell(w)$ is the length of w , which is

$$\ell(w) := \#\{k > i \mid w(k) < w(i)\}.$$

For example, here are \mathcal{M}_{463512} with $a_\bullet = 2 < 4$ and $\mathcal{M}_{3625714}$ for $a_\bullet = 2 < 5$

$$\begin{pmatrix} m_{1,1} & m_{1,2} & m_{1,3} & 1 & 0 & 0 \\ m_{2,1} & m_{2,2} & m_{2,3} & 0 & m_{2,5} & 1 \\ m_{3,1} & m_{3,2} & 1 & 0 & 0 & 0 \\ m_{4,1} & m_{4,2} & 0 & 0 & 1 & 0 \end{pmatrix} \quad \begin{pmatrix} m_{1,1} & m_{1,2} & 1 & 0 & 0 & 0 & 0 \\ m_{2,1} & m_{2,2} & 0 & m_{2,4} & m_{2,5} & 1 & 0 \\ m_{3,1} & 1 & 0 & 0 & 0 & 0 & 0 \\ m_{4,1} & 0 & 0 & m_{4,4} & 1 & 0 & 0 \\ m_{5,1} & 0 & 0 & m_{5,4} & 0 & 0 & 1 \end{pmatrix}$$

Recall that given a full rank $a_t \times n$ matrix M , $R_{a_\bullet}(M)$ is the flag of type a_\bullet whose a_i -dimensional linear subspace is the span of the first a_i rows of M . Under the map $M \mapsto R_{a_\bullet}(M)$, the set \mathcal{M}_w of partial echelon matrices parametrizes the Schubert cell $X_w^\circ E_\bullet$. This extends to any flag, not just E_\bullet .

Proposition 4.1. *Let Φ be a full rank $n \times n$ matrix. For $w \in W^{a_\bullet}$, the set \mathcal{M}_w of partial echelon matrices parametrizes the Schubert cell $X_w^\circ R_\bullet(\Phi)$ via the map $M \mapsto R_{a_\bullet}(M\Phi)$.*

Let Φ be a full rank $n \times n$ matrix and let M be a full rank $a_t \times n$ matrix. By the definition (4.1), for any $w \in W^{a_\bullet}$ the partial flag $R_{a_\bullet}(M)$ lies in $X_w R_\bullet(\Phi)$ if and only if, for all $i = 1, \dots, r$ and $j = 1, \dots, n$, we have that

$$(4.3) \quad \text{rank} \begin{pmatrix} M_{a_i} \\ \Phi_j \end{pmatrix} \leq a_i + j - \#\{k \leq a_i \mid w(k) \leq j\}.$$

Set $r_w(a_i, j) := a_i + j - \#\{k \leq a_i \mid w(k) \leq j\}$, which is this bound on rank. We use (4.3) to give the traditional determinantal equations defining a Schubert variety.

Proposition 4.2. *Let \mathcal{M} be a set of matrices parametrizing a subset Y of the flag manifold $\mathbb{F}\ell(a_\bullet, n)$, Φ a full-rank matrix, and $w \in W^{a_\bullet}$. The intersection of Y with the Schubert variety $X_w R_\bullet(\Phi)$ is defined in the coordinates \mathcal{M} for Y by the vanishing of the square $r_w(a_i, j)+1$ minors of $\begin{pmatrix} M_{a_i} \\ \Phi_j \end{pmatrix}$, for $j = 1, \dots, t$ and $k = 1, \dots, n$.*

The codimension of $X_w F_\bullet$ is $|w| := \dim(a_\bullet) - \ell(w)$. A *Schubert problem* on $\mathbb{F}\ell(a_\bullet, V)$ is a list w^1, \dots, w^s of permutations in W^{a_\bullet} such that $|w^1| + \dots + |w^s| = \dim(a_\bullet)$. By Kleiman's Transversality Theorem, if $F_\bullet^1, \dots, F_\bullet^s$ are general flags then the intersection

$$(4.4) \quad X_{w^1} F_\bullet^1 \cap X_{w^2} F_\bullet^2 \cap \dots \cap X_{w^s} F_\bullet^s$$

either is empty or is transverse and consists of finitely many points. As in Section 1, we may formulate a Schubert problem as a system of determinantal equations (given by Proposition 4.2) for membership in $X_{w^i} F_\bullet^i$ for $i = 1, \dots, s-1$ in the system \mathcal{M}_{w^s} of local coordinates for $X_{w^s} F_\bullet^s$. This will be overdetermined unless $|w^i| = 1$ for $i = 1, \dots, s-1$.

4.1. Primal-dual formulation for Schubert problems in $\mathbb{F}\ell(a_\bullet, V)$. Given a sequence $a_\bullet \subset [n-1]$, let a_\bullet^r be its reversal

$$a_\bullet^r : n-a_t < n-a_{t-1} < \dots < n-a_2 < n-a_1.$$

When $n = 9$, we have $(2, 3, 5)^r = (4, 6, 7)$. The annihilator map of Section 2 gives an isomorphism $\perp : \mathbb{F}\ell(a_\bullet, V) \xrightarrow{\sim} \mathbb{F}\ell(a_\bullet^r, V^*)$ in which $\perp(V_{a_\bullet})_{a_j^r} = (V_{a_{t+1-j}})^\perp$. Let $w_0 \in S_n$ be the permutation so that $w_0(i) = n+1-i$. Note that $\ell(w) = \ell(w_0 w w_0)$ and $w \in W^{a_\bullet}$ if and only if $w_0 w w_0 \in W^{a_\bullet^r}$. We give the version of Lemma 2.1 for $\mathbb{F}\ell(a_\bullet, V)$.

Lemma 4.3. *For a Schubert variety $X_w F_\bullet \subset \mathbb{F}\ell(a_\bullet, V)$, we have $\perp(X_w F_\bullet) = X_{w_0 w w_0} F_\bullet^\perp$.*

For $w \in W^{a_\bullet}$, define \mathcal{N}_w to be the set of $n \times (n-a_1)$ -matrices,

$$\mathcal{N}_w := \{(M \cdot w_0)^T \mid M \in \mathcal{M}_{w_0 w w_0}\},$$

which are obtained by first reversing the columns of a matrix in $\mathcal{M}_{w_0 w w_0}$ and then transposing it. We state the analog of Lemma 2.2 in this context.

Lemma 4.4. *Let Φ be a nonsingular $n \times n$ matrix. Under the map $N \mapsto C_{a_\bullet^r}(\Phi^{-1}N)$ the set \mathcal{N}_w of matrices for $w \in W^{a_\bullet}$ parametrizes the Schubert cell $X_{w_0 w w_0}^\circ(R_\bullet(\Phi))^\perp$ in $\mathbb{F}\ell(a_\bullet^r, V^*)$.*

Let us consider this duality in terms of bilinear equations in Stiefel coordinates. Suppose that M is a full rank $a_t \times n$ matrix with corresponding flag $V_{a_\bullet} := R_{a_\bullet}(M)$ and N a full rank $n \times (n-a_1)$ matrix with corresponding flag $U_{a_\bullet^r} := C_{a_\bullet^r}(N)$. As in Section 2, the columns of M correspond to a basis $\mathbf{e}_1, \dots, \mathbf{e}_n$ of V and the rows of N to the dual basis $\mathbf{e}_1^*, \dots, \mathbf{e}_n^*$ of V^* . Then the condition that $U_{a_\bullet^r} = V_{a_\bullet}^\perp$ is that

$$U_{a_{t+1-j}^r} = U_{n-a_j} = (V_{a_j})^\perp \quad \text{for } j = 1, \dots, t.$$

In terms of the matrices M and N this is

$$(4.5) \quad M_{a_j} \cdot N'_{n-a_j} = 0_{a_j \times (n-a_j)} \quad \text{for } j = 1, \dots, t.$$

(Here N'_{n-a_j} is the matrix formed by the first $n-a_j$ columns of N .) This set of $\sum_j a_j(n-a_j)$ bilinear equations is redundant. Write $\text{row}_i(M)$ for the i th row of the matrix M and $\text{col}_j(N)$ for the j th column of N . The following lemma is easily verified.

Lemma 4.5. *The set of bilinear equations (4.5) for $(M, N) \in \text{Mat}_{a_t \times n}(\mathbb{C}) \times \text{Mat}_{n \times (n-a_1)}(\mathbb{C})$ consists of the equations*

$$\text{row}_i(M) \cdot \text{col}_j(N) = 0,$$

for those (i, j) such that there is some k with $i \leq a_k$ and $j \leq (n-a_k)$. There are exactly $\dim a_\bullet$ such equations.

Write $MN = 0_{a_\bullet}$ for this set of $\dim a_\bullet$ bilinear equations. We give the primal-dual formulation of Schubert problems in $\mathbb{F}\ell(a_\bullet, V)$, which follows from the reduction to the dual diagonal as in Section 2.

Theorem 4.6. *Let (w^1, \dots, w^s) be a Schubert problem for $\mathbb{F}\ell(a_\bullet, V)$ and suppose that $F_\bullet^1, \dots, F_\bullet^s$ are general flags given by matrices Φ^1, \dots, Φ^s so that $F_\bullet^i = R_\bullet(\Phi^i)$. Then the system of equations*

$$(4.6) \quad (M\Phi^s)(\Phi^i)^{-1}N_i = 0_{a_\bullet} \quad \text{for } i = 1, \dots, s-1,$$

where $(M, N_1, \dots, N_{s-1}) \in \mathcal{M}_{w^s} \times \mathcal{N}_{w^1} \times \dots \times \mathcal{N}_{w^{s-1}}$ defines the instance of the Schubert problem (4.4) as the common zeroes of $(s-1) \cdot \dim(a_\bullet)$ bilinear equations in $(s-1) \cdot \dim(a_\bullet)$ variables, exhibiting the Schubert problem as a square system of equations.

As in Section 3, there are some improvements to this formulation. Rather than formulate a general result along the lines of Theorem 3.3, we will make a series of remarks indicating some reductions that are possible.

For $\alpha \in \binom{[n]}{a}$ with $a \in a_\bullet$, there are many permutations $w \in W^{a_\bullet}$ with $w|_a = \alpha$. (These are in bijection with $W_{\{a\}}/W_{a_\bullet} = (S_a \times S_{n-a})/(S_{a_1} \times S_{a_2-a_1} \times \dots \times S_{n-a_t})$.) Write $w(\alpha, a_\bullet)$, or $w(\alpha)$ when a_\bullet is assumed, for the longest such permutation in W^{a_\bullet} . For example, if $n = 8$ and $a_\bullet = (2 < 4 < 5)$, then $w(2, 5, 6, 8) = 68.25.7.134$. Observe that $|\alpha| = |w(\alpha)|$.

Schubert varieties given by permutations $w(\alpha)$ have a form which we exploit to reduce the number of equations and variables in (4.6). By (4.2), if $\alpha \in \binom{[n]}{a}$ where $a \in a_\bullet$, then

$$X_{w(\alpha, a_\bullet)}F_\bullet = \{V_{a_\bullet} \in \mathbb{F}\ell(a_\bullet, V) \mid V_a \in X_\alpha F_\bullet\}.$$

In particular, we have

$$V_{a_\bullet} \in X_{w(\alpha, a_\bullet)}F_\bullet \iff (V_a)^\perp \in X_{\alpha^\perp}F_\bullet^\perp.$$

Remark 4.7. Suppose that one permutation in Theorem 4.6, say w^i , has the form $w(\alpha, a_\bullet)$ for $\alpha \in \binom{[n]}{a}$ with $a \in a_\bullet$. Then we may replace the factor \mathcal{N}_{w^i} by the space of $n \times (n-a)$ -matrices \mathcal{N}_α and the i th set of equations in (4.6) by

$$(M_a \Phi^s)(\Phi^i)^{-1}N = 0_{a \times (n-a)},$$

where $N \in \mathcal{N}_\alpha$. This uses $\dim a_\bullet - a(n-a)$ fewer equations and variables.

Remark 4.8. Suppose that two of the permutations have this form, say $w^i = w(\alpha, a_\bullet)$ and $w^j = w(\beta, a_\bullet)$ for $\alpha, \beta \in \binom{[n]}{a}$ with $a \in a_\bullet$. There is a matrix Φ so that $F_\bullet^i = R_\bullet(\Phi)$ and $F_\bullet^j = R_\bullet(w_0\Phi)$ are the pair of opposite flags associated to Φ . In Theorem 4.6 we can replace the factors $\mathcal{N}_{w^i} \times \mathcal{N}_{w^j}$ by the factor \mathcal{N}_α^β and the equations indexed by i and j in (4.6) by the equations

$$(M_a \Phi^s) \cdot (\Phi^{-1}N) = 0_{a \times (n-a)},$$

where $N \in \mathcal{N}_\alpha^\beta$. This uses $2 \dim a_\bullet - a(n-a)$ fewer equations and variables.

Remark 4.9. If some permutation w^i has $|w^i| = 1$, then it has the form $w(\alpha, a_\bullet)$ where $\alpha \in \binom{[n]}{a}$ satisfies $|\alpha| = 1$ so that $\alpha = \square$. Write $w(\square, a, a_\bullet)$ or $w(\square, a)$ for this permutation. We may remove the factor \mathcal{N}_{w^i} and replace the i th set of equations in (4.6) with the single determinantal equation

$$\det \begin{pmatrix} M_a F^s \\ \Phi_{n-a}^i \end{pmatrix} = 0.$$

This uses $\dim a_\bullet - 1$ fewer equations and variables.

Remark 4.10. We generalize Remark 4.7. Let $b_\bullet \subsetneq a_\bullet$ and suppose that $v \in W^{b_\bullet}$. Write $w(v, a_\bullet) \in W^{a_\bullet}$ for the longest permutation in W^{a_\bullet} that lies in the coset vW_{b_\bullet} . Then $|w(v, a_\bullet)| = |v|$, or more precisely $\dim a_\bullet - \ell(w(v, a_\bullet)) = \dim b_\bullet - \ell(v)$. For example, if $n = 8$, $a_\bullet = (2 < 4 < 5)$ and $v = 68257134 \in W^{(2 < 5)}$, then $w(v, a_\bullet) = 68.57.2.134$. Also, if $v = 47681235 \in W^{(2 < 4)}$, then $w(v, a_\bullet) = 47.68.5.123$, and if $v = 34685127 \in W^{(4 < 5)}$, then $w(v, a_\bullet) = 68.34.5.127$.

A flag $V_{a_\bullet} \in X_{w(v, a_\bullet)} F_\bullet$ if and only if the subflag V_{b_\bullet} consisting of the subspaces with dimensions in b_\bullet lies in the Schubert variety $X_v F_\bullet$ of $\mathbb{F}\ell(b_\bullet, V)$. In this case we have that

$$V_{a_\bullet} \in X_{w(v, a_\bullet)} F_\bullet \iff (V_{b_\bullet})^\perp \in X_{w_0 v w_0} F_\bullet^\perp \subset \mathbb{F}\ell(b_\bullet^r, V^*).$$

If one of the permutations w^i in Theorem 4.6 has the form $w(v, a_\bullet)$ for $v \in W^{b_\bullet}$, then we may replace the set \mathcal{N}_{w^i} of matrices parametrizing $X_{w_0 w^i w_0} F_\bullet^\perp$ by the set of matrices \mathcal{N}_v parametrizing $X_{w_0 v w_0} F_\bullet^\perp \subset \mathbb{F}\ell(b_\bullet^r, V^*)$, and the i th set of equations in (4.6) by

$$(M \Phi^s)(\Phi^i)^{-1}N = 0_{b_\bullet},$$

for $N \in \mathcal{N}_v$. This uses $\dim a_\bullet - \dim b_\bullet$ fewer equations and variables.

Example 4.11. Suppose $n = 8$ and $a_\bullet = (2 < 4 < 5)$. Note that

$$\dim a_\bullet = 2(8-2) + (4-2)(8-4) + (5-4)(8-5) = 23.$$

Using $(w)^m$ to indicate that w is repeated m times, consider the Schubert problem,

$$48573126, (78453126)^2, (68574123)^2, (78465123)^3, 47385126,$$

with 128 solutions. The first permutation is $w(v)$ where $v = 48.357.126 \in W^{(2 < 5)}$, the next two are $w(\alpha)$ where $\alpha = (3, 4, 6, 7, 8)$, the next two are $w(\square, 2)$, and the next three are $w(\square, 4)$. Only the last permutation does not reduce to a smaller flag manifold.

By Theorem 4.6, this has a formulation as a square system of $8 \cdot 23 = 184$ bilinear equations and variables.

However, we may use the reduction of Remark 4.10 with the first permutation to use $2 = \dim a_{\bullet} - \dim b_{\bullet}$ fewer equations, then the reduction of Remark 4.8 with the second and third permutations to use $2 \cdot \dim a_{\bullet} - 5(8 - 5) = 31$ fewer equations, and then the reduction of Remark 4.9 on the next five permutations to use $5(23-1) = 110$ fewer equations, and thereby obtain a square system with $184 - 2 - 31 - 110 = 41$ variables, 36 bilinear equations, and 5 determinantal equations involving 2 quadratic determinants and 3 quartic determinants.

We used Bertini [2] to solve a random real instance of this Schubert problem with the square formulation of 41 variables. This consumed about 2.95 gigaHertz-days of processing power to calculate 128 approximate solutions, and the output suggests that 42 solutions are real flags while the other 86 are nonreal. As in Example 3.5, this computation may be done in parallel efficiently.

We used rational arithmetic in alphaCertified [13] to prove that the 128 solutions given by Bertini are approximate solutions to the computed instance, and that they correspond to distinct solutions. As in Example 3.5, available theorems guarantee that we found approximations for every solution to the instance. We also verified that 42 of the approximate solutions correspond to real solutions. This calculation took 1.78 gigaHertz-hours of processing time. This computation may also be done efficiently in parallel. Details for these computations and those in Example 3.5 may be found at http://www.unk.edu/academics/math/_files/primal-dual.html.

REFERENCES

- [1] D. Armentano, *Complexity of path-following methods for the eigenvalue problem*, Found. Comput. Math. **14** (2014), no. 2, 185–236.
- [2] Daniel J. Bates, Jonathan D. Hauenstein, Andrew J. Sommese, and Charles W. Wampler, *Bertini: Software for numerical algebraic geometry*, Available at bertini.nd.edu with permanent doi: [dx.doi.org/10.7274/R0H41PB5](https://doi.org/10.7274/R0H41PB5).
- [3] C. Beltrán and A. Leykin, *Certified numerical homotopy tracking*, Exp. Math. **21** (2012), no. 1, 69–83.
- [4] C. Beltrán and A. Leykin, *Robust certified numerical homotopy tracking*, Found. Comput. Math. **13** (2013), no. 2, 253–295.
- [5] L. Blum, F. Cucker, M. Shub, and S. Smale, *Complexity and real computation*, Springer-Verlag, New York, 1998, With a foreword by Richard M. Karp.
- [6] J.-P. Dedieu and M. Shub, *Multihomogeneous newton methods*, Math. Comp. **69** (2000), no. 231, 1071–1098.
- [7] A. Eremenko and A. Gabrielov, *Rational functions with real critical points and the B. and M. Shapiro conjecture in real enumerative geometry*, Ann. of Math. (2) **155** (2002), no. 1, 105–129.
- [8] W. Fulton, *Young tableaux*, London Mathematical Society Students Texts, 35, Cambridge University Press, Cambridge, 1997.
- [9] L. García-Puente, N. Hein, C. Hillar, A. Martín del Campo, J. Ruffo, F. Sottile, and Z. Teitler, *The Secant Conjecture in the real Schubert Calculus*, Exper. Math. **21** (2012), no. 3, 252–265.
- [10] J.D. Hauenstein, I. Haywood, and A.C. Liddell, Jr., *An a posteriori certification algorithm for Newton homotopies*, Proc. ISSAC 2014, to appear, 2014.
- [11] J.D. Hauenstein, N. Hein, C. Hillar, A. Martín del Campo, Frank Sottile, and Zach Teitler, *The Monotone Secant Conjecture in the real Schubert calculus*, MEGA11, Stockholm, 2011.
- [12] J.D. Hauenstein and F. Sottile, *Algorithm 921: alphaCertified: Certifying solutions to polynomial systems*, ACM Trans. Math. Softw. **38** (2012), no. 4, 28.

- [13] Jonathan D. Hauenstein and Frank Sottile, *alphaCertified: Software for certifying numerical solutions to polynomial equations*, Available at www.math.tamu.edu/~sottile/research/stories/alphaCertified.
- [14] N. Hein, C. Hillar, and F. Sottile, *Lower bounds in real Schubert calculus*, São Paulo Journal of Mathematics **7** (2013), no. 1, 33–58.
- [15] B. Huber, F. Sottile, and B. Sturmfels, *Numerical Schubert calculus*, J. Symb. Comp. **26** (1998), no. 6, 767–788.
- [16] S.L. Kleiman, *The transversality of a general translate*, Compositio Math. **28** (1974), 287–297.
- [17] A. Leykin and F. Sottile, *Galois groups of Schubert problems via homotopy computation*, Math. Comp. **78** (2009), no. 267, 1749–1765.
- [18] Abraham Martín del Campo and Frank Sottile, *Experimentation in the Schubert calculus*, [arXiv.org/1308.3284](https://arxiv.org/abs/1308.3284), 2013.
- [19] E. Mukhin and V. Tarasov, *Lower bounds for numbers of real solutions in problems of Schubert calculus*, 2014, [arXiv:1404.7194](https://arxiv.org/abs/1404.7194).
- [20] E. Mukhin, V. Tarasov, and A. Varchenko, *The B. and M. Shapiro conjecture in real algebraic geometry and the Bethe ansatz*, Ann. of Math. (2) **170** (2009), no. 2, 863–881.
- [21] E. Mukhin, V. Tarasov, and A. Varchenko, *Schubert calculus and representations of the general linear group*, J. Amer. Math. Soc. **22** (2009), no. 4, 909–940.
- [22] J. Ruffo, Y. Sivan, E. Soprunova, and F. Sottile, *Experimentation and conjectures in the real Schubert calculus for flag manifolds*, Exper. Math. **15** (2006), no. 2, 199–221.
- [23] H. Schubert, *Kalkül der abzählenden Geometrie*, Springer-Verlag, 1879, reprinted with an introduction by S. Kleiman, 1979.
- [24] H. Schubert, *Anzahl-Bestimmungen für lineare Räume beliebiger Dimension*, Acta. Math. **8** (1886), 97–118.
- [25] H. Schubert, *Die n-dimensionalen Verallgemeinerungen der fundamentalen Anzahlen unseres Raums*, Math. Ann. **26** (1886), 26–51, (dated 1884).
- [26] H. Schubert, *Losung des Charakteritiken-Problems für lineare Räume beliebiger Dimension*, Mittheil. Math. Ges. Hamburg (1886), 135–155, (dated 1885).
- [27] M. Shub and S. Smale, *Complexity of Bézout’s theorem. I. Geometric aspects*, J. Amer. Math. Soc. **6** (1993), no. 2, 459–501.
- [28] S. Smale, *Newton’s method estimates from data at one point*, The merging of disciplines: new directions in pure, applied, and computational mathematics, Springer, New York, 1986, pp. 185–196.
- [29] A.J. Sommese and C. W. Wampler, II, *The numerical solution of systems of polynomials*, World Scientific Publishing Co. Pte. Ltd., Hackensack, NJ, 2005.
- [30] F. Sottile, *Real Schubert calculus: Polynomial systems and a conjecture of Shapiro and Shapiro*, Exper. Math. **9** (2000), 161–182.
- [31] F. Sottile, *Frontiers of reality in Schubert calculus*, Bull. Amer. Math. Soc. (N.S.) **47** (2010), no. 1, 31–71.
- [32] F. Sottile, *Real solutions to equations from geometry*, University Lecture Series, vol. 57, American Mathematical Society, 2011.
- [33] F. Sottile, R. Vakil, and J. Verschelde, *Solving Schubert problems with Littlewood-Richardson homotopies*, Proc. ISSAC 2010 (Stephen M. Watt, ed.), ACM, 2010, pp. 179–186.

JONATHAN D. HAUENSTEIN, DEPARTMENT OF APPLIED AND COMPUTATIONAL MATHEMATICS AND STATISTICS, UNIVERSITY OF NOTRE DAME, NOTRE DAME, INDIANA 46556, USA

E-mail address: hauenstein@nd.edu

URL: <http://www.nd.edu/~jhauenst>

NICKOLAS HEIN, DEPARTMENT OF MATHEMATICS, UNIVERSITY OF NEBRASKA AT KEARNEY, KEARNEY, NEBRASKA 68849, USA

E-mail address: heinnj@unk.edu

URL: http://www.unk.edu/academics/math/faculty/About_Nickolas_Hein/

FRANK SOTTILE, DEPARTMENT OF MATHEMATICS, TEXAS A&M UNIVERSITY, COLLEGE STATION,
TEXAS 77843, USA

E-mail address: `sottile@math.tamu.edu`

URL: `http://www.math.tamu.edu/~sottile/`