# CSE 30151 Theory of Computing: Homework 4 GNFAs, Regex, Pumping Lemma 

Version 1: Sept. 14, 2017

## Instructions

- Unless otherwise specified, all problems from "the book" are from Version 3. When a problem in the International Edition is different from Version 3, the problem will be listed as V3:x.yy/IE:x.zz, where $\mathrm{x} . \mathrm{zz}$ is the equivalent number. When Version 2 has a different number, it will be listed as V3:x.yy/V2:x.zz. If either IE or V2 does not have a matching number, the problem text will be duplicated.
- You can prepare your solutions however you like (handwriting, $\mathrm{A}_{\mathrm{E}} \mathrm{X} \mathrm{X}$, etc.), but you must submit them in PDF. You can scan written solutions in the library or using a smartphone (with a scanner app like CamScanner). It is up to you to ensure that submissions are legible.
- Please give every PDF file a unique filename.
- If you're making a complete submission (all problems), name your PDF file netid-hw2.pdf, where netid is replaced with your NetID.
- If you're submitting some problems now and other problems later, name your file netid-hw2-123.pdf, where 123 is replaced with just the problems you are submitting now.
- If you use the same filename twice, only the most recent version will be graded.
- The time of submission is the time the most recent file was uploaded.
- If you use $\mathrm{IA}_{\mathrm{E}} \mathrm{X}$ and want to draw something like a state diagram, consider using the tikz package. A reference document is on the website under "Assignments".
- Submit your PDF file in Sakai. Don't forget to click the Submit (or Resubmit) button!


## Practice Problems

These problems are from the book, and most have solutions listed for them. They are listed here for you to practice on as needed and any answers you generate should not be submitted. You are free to discuss these with others, but you are not allowed to post solutions to any public forum.

1. 1.10 Regex to NFA
2. 1.20 Regex to strings
3. 1.21a DFA to regex
4. 1.28 regex to NFA
5. $1.29 \mathrm{a}, \mathrm{c}$ Pumping Lemma to show non-regularity.

## Book Exercises

These problems are found in the text book and are to be answered and submitted by each student. You are to solve them yourself. Use of solution manuals from any source or shared solutions is a violation of the ND Honor Code. You are also not allowed to show your solutions to another student.

1. (5 points) 1.12 Language to DFA and regular expression. Hint: what does not including "ab" mean to any strings in the language.
2. (5 points) 1.19a Regular expression to NFA
3. (5 points) 1.21 b DFA to regex. Show steps
4. (5 points) 1.29b Pumping Lemma
5. ( 5 points) 1.37 Show a language is regular. (Hint: remember that if for the binary number $\mathrm{x}, \bmod (x, n)=$ $k$, then the number $\bmod (2 x+b, n)$, where b is another binary digit, is the same as $\bmod (2 k+b, n))$

## Non-book Problems

The following problems are not found in the text book. You are to solve them yourself. Use of any resource you used other than the text book or class notes must be cited. You are also not allowed to show your solutions to another student.
6. (5 points) For the following NFA ompute the regular expression accepted by converting it to a GNFA and then using the CONVERT algorithm of from Lemma 1.60 on it. Show work. Feel free to use variables like $R_{i}$ to stand for intermediate expressions on an edge (with $R_{i}$ written out below), but show entire final answer. You don't need to simplify.

7. (5 points) Do the same as above for the following NFA, i.e. compute the regex.

8. (5 points) Show using the Pumping Lemma that the language $L=\left\{(a b)^{n}(c d)^{n} \mid n \geq 0\right\}$ is not regular.
9. (5 points) Show that the language generated from $\Sigma=\{a, b\}$ where the only constraint is that the number of a's in any string is less than the number of b's, is not regular.
10. (5 points) A string homomorphism h is a function that takes each character in its input string some alphabet $\Sigma^{*}$, replaces it by a string (typically from a different alphabet $\Sigma^{\prime}$ ), and concatenates all resulting substrings together, i.e. $h: \Sigma^{*} \rightarrow\left(\Sigma^{\prime}\right)^{*}$. As an example if $h(0)=a b$ and $h(1)=\varepsilon$, then $h(00110)=a b a b a b$
Prove that if L is a regular language over $\Sigma$, then the language $L^{\prime}=\{h(w) \mid w$ is in $L\}$ is a regular language over $\Sigma^{\prime}$. Hint: show by induction over the construction of regular expressions.

