

Depth-First Search and Its Use Case in Distributed Systems Debugging

Nate Kremer-Herman



Task may create files, interact with and set environment variables, etc.

Project goal

- Trace and traverse the causal history of events in a distributed system.
- Use depth first search as a querying tool for finding causal history.
- Find the minimum set of events needed in the log to build a proper historical trace (out of the scope of this project).

Implementation techniques

Perl

- Implemented Work Queue master in Perl.
 - ▷ Could also use C or Python.
 - ▷ Distributes each *error* tagged node.
- Workers use the iterative algorithm.
 - ▷ Can only use recursion to a certain depth.


```
my $wq = Work_Queue->new( port => 0, name => "traverse", catalog => 1);
$wq->specify catalog server("catalog.cse.nd.edu", 9097);
my sport = $wq->port();
print(STDERR "Work Queue listening on port $port...\n");
                                                           ry 8192 -- disk 1024 $workers > /dev/null");
system ("condor submit
my $i = 0;
while($i <= $runs) {</pre>
   my $jobs = 0;
   my $tool = "subtraverse";
   my $traversed = 0;
   my Gouts;
   my $epoch = time();
   foreach my Se (Gerrors) {
       my $out = "traversal.$jobs.out";
       my $command = "perl $tool -i $input -r $e > $out";
       my st = Work Queue::Task->new($command);
       $t->specify input file(local name => $tool, remote name => $tool);
       $t->specify input file(local name => $input, remote name => $input);
       $t->specify output file($out);
       $wq->submit($t);
       push(@outs, Sout);
       $jobs++;
   print(STDERR "All tasks submitted.\n");
   while(!$wq->empty()) {
       my $t = $wq->wait(10);
       if($t) {
          $jobs--;
```

Sequential notional summary (for worker nodes)

- 1 **procedure** DFS-iterative(*G*,*v*):
- 2 push v on a stack, S
- 4 **while** *S* is not empty
 - *v* = *S*.pop()

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- 6 **if** *v* has not been visited in this round:
 - label v as visited
- 8 for all child edges of v do
- 9 if child has a matching attribute with v: //file or environment variable
 10 S.push(child)

Updated complexity analysis

- Time complexity is still O(|V| + |E|)
 - ▷ Worst case, we look at all vertices.
 - ▷ Best case, we look at no vertices (no errors!).
- Space is now O(W (|V| + |E|))
 - ▷ Where W is the number of workers.
 - Graph must be sent to each worker once, then cached for future tasks.

Datasets

- All datasets are synthetic
 - Each is a binary graph
 - ▷ Ran out of time to produce greater variation.
 - Generated via Perl script
- Number of nodes ranges from 10 1,000,000
 - Realistic dataset size O(100) O(10,000)
 - ▷ Tiny: 10 nodes
 - ⊳ Small: 100 nodes
 - \triangleright

...

Colossal: 1,000,000 nodes

Serial Traversal Time for Varying Graph Size

Size of Graph (Nodes)

Parallel Traversal Time for Varying Graph Size

Size of Graph (Nodes)

Peak Provisioning of a Parallel Application

Scale

Conclusions

- There exists some scale where parallel is better than serial traversal.
 - ▷ Did not find that within realistic data sizes.
- I now have the graph traversal backend for my future research.
 - ▷ Need to make a graphifier for debug logs.
- Do not make DFS parallel.
 - ▷ Wonder if Gremlin can be used instead.

Questions?

