

Lemma: Let γ, γ' be two cell decompositions of a cpt 2-mfd Σ .

Then $\chi(\Sigma, \gamma) = \chi(\Sigma, \gamma')$

Argument: • by moving Γ, Γ' a bit, can make their vertex sets disjoint and have fin. many intersections between edges of Γ and Γ'

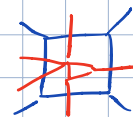
• can construct γ'' - a common refinement of γ, γ' (i.e. γ'' is obtained from γ, γ' by inductively adding new vertices on edges \downarrow or adding new edges btw vertices of a polygon



how? $V_{\gamma''} = V_{\gamma} \sqcup V_{\gamma'} \sqcup (\Gamma \cap \Gamma')$

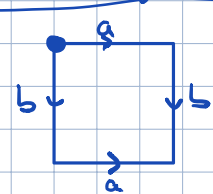
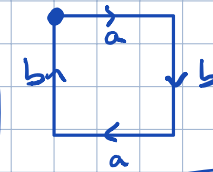

intersection pts of edges of Γ and Γ'

edges of $\gamma'' =$ segments of edges of γ, γ' connecting vertices of γ''



A combinatorial description of compact connected 2-manifolds

2-mfds obtained by gluing edges of a single polygon \longleftrightarrow "words"

2-mfd	edge gluing	word
torus		$aba^{-1}b^{-1}$
Klein bottle		$aba^{-1}b$
$\mathbb{R}P^2$		aa

surface given by edge-gluing of an n-gon
(edges \rightarrow labels)
orientation

word: fix a vertex and go clockwise from it

$$W = x_1 x_2 \dots x_n$$

$$x_i = (e_i)^{\pm 1}$$

label of i^{th} edge e_i

power ± 1 : $+1$ if e_i is oriented clockwise
 -1 if counterclockwise

(seq. of "letters" $\in A = \{a, a^{-1}, b, b^{-1}, \dots\}$; $a, b, \dots \in L$ set of labels)

reverse direction:

$$\text{word } W = l_i^{\pm 1} \dots l_n^{\pm 1}$$

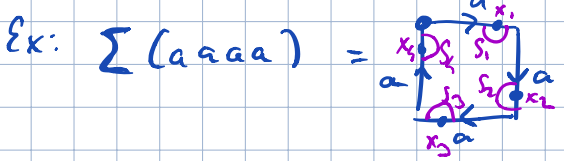
$$\mathbb{P}_n / \sim_W =: \Sigma(W)$$

n-gon with a distinguished vertex

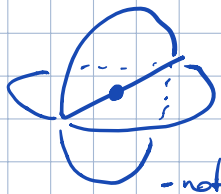


mark i^{th} edge e_i (clockwise) by l_i , orient clockwise if l_i^{+1}
counterclockwise if l_i^{-1} appears in W .

Rem $\Sigma(W)$ is not always a manifold.



open nbhd of $[x_i] = S_1 \cup \dots \cup S_k / \sim_W$



- gluing of \leq semi-disks.

- not a manifold!

Lemma ^A 1) let W be a word built from labels in a set L . Let $L \leftrightarrow L'$ be a bijection of sets and let W' be the word obtained by replacing each occurrence of $l \in L$ with $l' \in L'$ where l' corresponds to l via bijection. Then

$$\Sigma(W) \approx \Sigma(W')$$

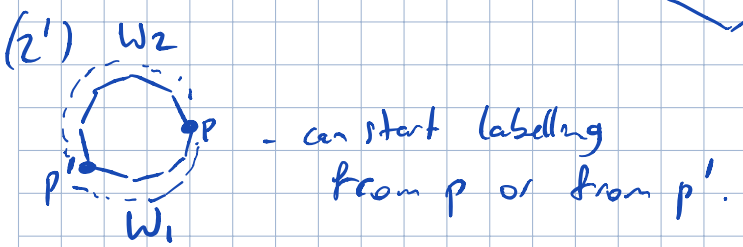
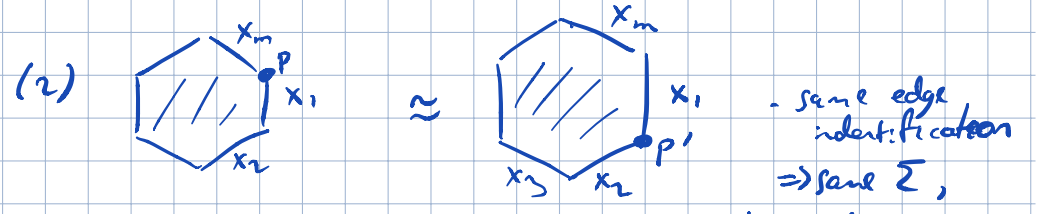
$$(2) \Sigma(x_1 x_2 \dots x_m) \approx \Sigma(x_2 x_3 \dots x_m x_1)$$

more generally: for W_1, W_2 two words with letters in A ,

$$(\textcircled{a}) \Sigma(W_1 W_2) \approx \Sigma(W_2 W_1)$$

$x_1 \dots x_m \quad y_1 \dots y_n$
concatenation
of words

Proof: (1): obvious

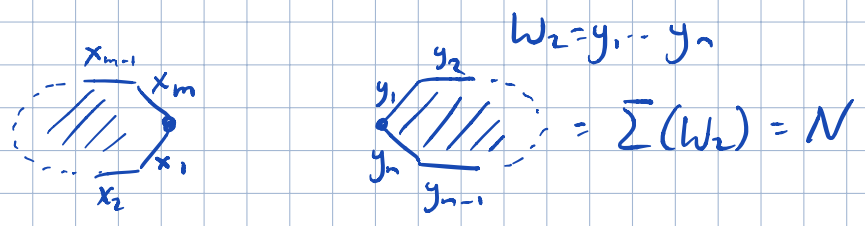


Proposition Let M, N be two cpt, conn 2-mpds, $M = \Sigma(W_1)$, $N = \Sigma(W_2)$, W_1, W_2 - words from disjoint alphabets. Then the connected sum is

$$M \# N = \Sigma(W_1 W_2)$$

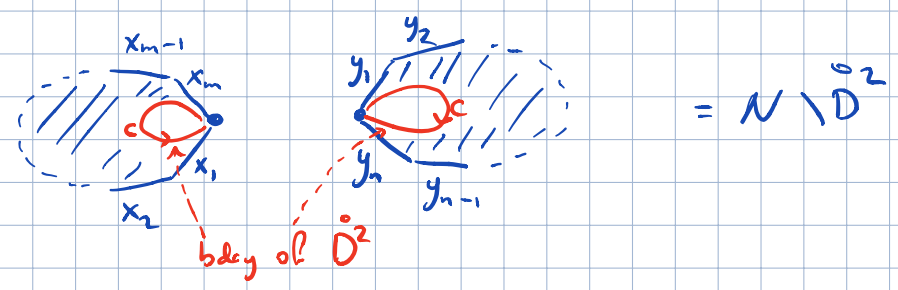
Proof $W_1 = x_1 \dots x_m$

$$M = \Sigma(W_1) =$$



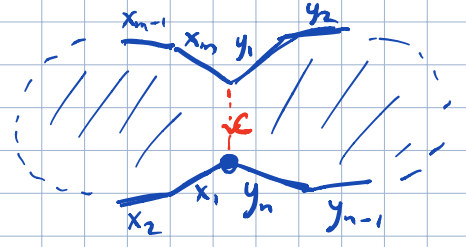
(remove $\overset{\circ}{D}^2$ from M and N)

$$M \setminus \overset{\circ}{D}^2 =$$



(glue along the circle c)

$$M \# N =$$



$$\approx \Rightarrow M \# N = \sum (x_1 \dots x_m y_1 \dots y_n) = \sum (W_1 W_2) \quad \square$$

(4)

Corollary: (1) $\Sigma_g = \underbrace{T \# \dots \# T}_g \approx \sum (a_1 b_1 a_1^{-1} b_1^{-1} a_2 b_2 a_2^{-1} b_2^{-1} \dots a_g b_g a_g^{-1} b_g^{-1})$
 $\approx \sum (a_1 b_1 a_1^{-1} b_1^{-1}) \# \dots \# \sum (a_g b_g a_g^{-1} b_g^{-1})$

(2) $X_k = \underbrace{\mathbb{R}P^2 \# \dots \# \mathbb{R}P^2}_k \approx \sum (a_1 a_1 a_2 a_2 \dots a_k a_k)$
 $\approx \sum (a_1 a_1) \# \dots \# \sum (a_k a_k)$

Proposition^B: Let W_1, W_2, W_3 be words and a a letter not occurring in them.

Then there are homeomorphisms

(*) $\Sigma(W_1 a W_2 a W_3) \approx \Sigma(W_1 a a W_2^{-1} W_3)$,

(**) $\Sigma(W_1 a W_2 a W_3) \approx \Sigma(W_1 W_2^{-1} a a W_3)$

where W_2^{-1} is the "inverse" of the word W_2 ; $W_2 = x_1 \dots x_n \rightarrow W_2^{-1} = x_n^{-1} \dots x_1^{-1}$.
 (as for a group product)

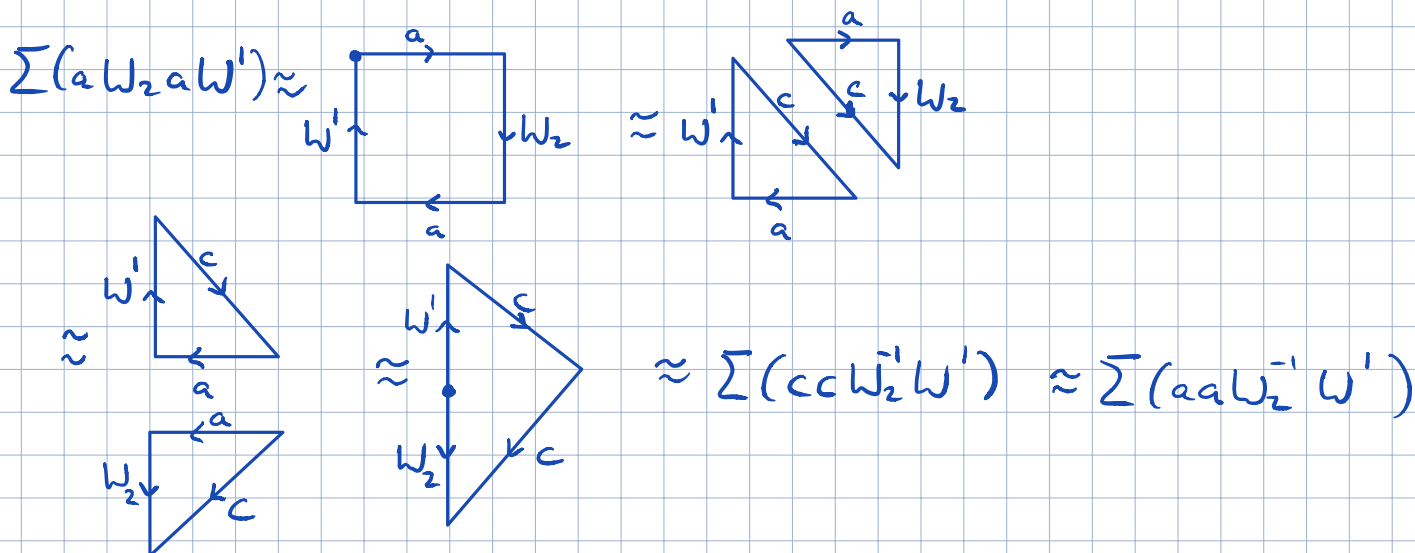
Proof: <let's check (*)>

$$\Sigma(W_1 a W_2 a W_3) \approx \Sigma(a W_2 a W_1')$$

(@) $\underbrace{W_3 W_1}_{W_1' W_3}$

$$\Sigma(W_1 a a W_2^{-1} W_3) \approx \Sigma(a a W_2^{-1} W_1')$$

So, we want to prove: $\Sigma(a W_2 a W_1') \approx \Sigma(a a W_2^{-1} W_1')$



(**) is similar - we cut the square by the other diagonal

□

Application: proof of $T \# \mathbb{R}P^2 \approx \mathbb{R}P^2 \# \mathbb{R}P^2 \# \mathbb{R}P^2$

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Indeed: $T \# \mathbb{R}P^2 = \Sigma (aba^{-1}b^{-1}) \# \Sigma (cc) \approx \Sigma (aba^{-1}b^{-1}cc)$

$\approx \Sigma (abc \underline{bac}) \approx \Sigma (ab \underline{bc^{-1}ac}) \approx \Sigma (bb \underline{c^{-1}aca})$

$\approx \Sigma (bb \underline{c^{-1}c^{-1}aa}) \approx \Sigma (bb) \# \Sigma (c^{-1}c^{-1}) \# \Sigma (aa) \approx \mathbb{R}P^2 \# \mathbb{R}P^2 \# \mathbb{R}P^2$

□

• Idea of proof of "constructive part" of the classification THM
<that any cpt conn 2-nd is $\approx \Sigma_g$ or X_k >

(1) show that every Σ admits a triangulation (cell decomp where polygons = triangles)
→ give each edge its own label

⇒ $\Sigma = \coprod \text{polygons}$
gluing edges carrying same label

(2) Reduce the number of polygons by one by gluing a pair of edges with same label belonging to different polygons

→ inductively, reduce to a single polygon

(3) Use the moves of Lemma A, Prop. B to show that labeling of edges of the polygon can be modified, without changing the homeo type of the resulting quotient, to obtain the stand. labeling for Σ_g or X_k .